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SIMCOMP VERSION 3.0 USER'S MANUAL

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TABLE OF CONTENTS

List of Examples

-v-

ABSTRACT

SIMCOMP Version 3.0 is a FORTRAN-like computer simulation language designed to ease the development and implementation of compartmental flow/ discrete event simulations. The language is an extension and refinement of SIMCOMP Version 2.0. The system is designed to reduce the programming overhead required while sufficient flexibility to solve certain problems. Compartmental-flow simulations are defined by specifying the flow rates between compartments and may be in either difference or differential equation form. Discrete events can be included in the form of event routines which are controlled by a dynamic event scheduler. Various forms of tabular and graphic output can be easily requested. An execution time interpretive debugging facility is also included. The syntactical rules for writing SIMCOMP programs are presented in this document along with a number of examples.

Introduction

The development of a second major simulation compiler, SIMCOMP Version 3.0, perhaps requires some justification. SIMCOMP Version 3.0 represents a series of refinements which have extended the capabilities of SIMCOMP Version 2.0 (Gustafson and Innis 1972). The refinements which have been incorporated into the Version 3.0 compiler are best understood in three broad categories. These categories are (i) refinements and extensions due to conceptual advances in modeling paradigms; (ii) refinements which ease the coding of SIMCOMP source programs; and (iii) the development of more efficient algorithms within the compiler and in the generated object program.

As a means of abstraction and representation of ecological systems, computer simulation models have taken many forms. Since many successful ecological simulations have adopted the state equation approach, we do not wish to abandon this paradigm. Instead, SIMCOMP Version 2.0 was designed around an extension of the state equation paradigm, namely the conceptualization of an ecological system as a set of compartmental-flow equations. SIMCOMP Version 3.0 retains the capability to represent systems as sets of compartmental-flow equations. The program organizes these equations into difference equations and provides a solution algorithm.

Certain categories of ecological phenomena appear strained when described in terms of state equations. These categories include (i) phenomena which admit to a heuristic description; (ii) phenomena with a low probability of occurrence; and (iii) phenomena which are best described as stochastic processes. The above categories are not meant to be exhaustive or mutually exclusive. The feature incorporated into SIMCOMP Version 3.0 to facilitate the modeling of such phenomena includes event routines and an event scheduler. In addition, a number of distribution functions for the generation of stochastic

variates are available. Event and compartmental-flow simulations may be combined in the same simulation.

Changes in the format of SIMCOMP Version 3.0 statements are in large measure due to the comments and suggestions of the users of Version 2.0. Source program statements were simplified when possible. A number of new statements have increased the capabilities of the language. For example, the fact has been recognized that many times in large ecosystem simulations flow equations are developed which take similar mathematical form, varying only in the values of the parameters. SIMCOMP Version 3.0 provides for such cases with the capability to iteratively declare flows. Additionally, an execution time interpretive variable dump facility has been provided to ease the debugging process.

The execution characteristics of the compiler and the generated object code have been significantly improved in Version 3.0. Execution times and core requirements have been reduced. A complete description of the system and its operation is contained in the "SIMCOMP Version 3.0 Maintainance Document" (Stevens and Gustafson 1973).

1. SIMCOMP SIMULATIONS.

Every simulation language is designed to ease the translation of real-world phenomena into computer simulation models. Each language is designed to model a particular class of real systems. Some notable examples are the application of MIMIC (Control Data Corporation 1972) to the simulation of continuous physical systems and the application of SIMSCRIPT (Markowitz et al. 1963) to the simulation of queuing and inventory systems. SIMCOMP was designed primarily to model ecological systems. As is the case with other simulation languages, the realm of applicability of SIMCOMP is certainly not limited to the field for which it was designed (ecology).

The principle features of SIMCOMP were designed with the modeling of ecological systems in mind. This desire required some broad generalizations about the nature of ecological systems. These generalizations include the following:

(1) Ecological systems can be viewed, in part, as systems of continuously varying variables.

(2) Ecological systems can be visualized as a system of compartments linked by material or energy transfers or flows.

(3) The flow of material or energy depends on other states and driving variables of the system (information flows).

(4) Ecological systems contain components which can be visualized as discrete-valued variables.

(5) Ecological systems contain processes which can be visualized as events occurring discretely through time.

SIMCOMP is designed to model phenomena in the above categories. Sections 1.1 and 1.2 are presented to formally define the paradigms employed by SIMCOMP to implement these generalizations. Sections 2 and 3 describe the syntax and coding procedures for writing SIMCOMP simulations.

1.1 Flow-Oriented Continuous Simulations.

A broad variety of techniques have been developed to model and simulate many systems. Differential equations and difference equations have been most used in the development of ecological models. As greater reality and resultant complexity are introduced, the solutions have become more intractable and computers have been employed. Computer simulation in this context has come to mean the numerical solution of a simultaneous set of differential or difference equations.

Although many solution schemes are available, one of the most versatile approaches is to view the simulation as an initial value problem. The initial value problem for first-order difference equations takes the following form. Let the amount of material or energy in the ith compartment at time t be represented by $x_i(t)$. For a system of n compartments, the state of the system at any time t can be expressed as a vector

$$x(t) = \left(x_1(t), x_2(t), \ldots, x_n(t) \right)$$

Let a change in the state of the system over some time interval, say Δt from time t to time $t + \Delta t$, be represented by

$$\Delta x(t) = \left(\Delta x_1(t), \Delta x_2(t), \ldots, \Delta x_n(t) \right) .$$

In general, the $\Delta x_i(t)$ are functions which may depend upon,

(1) the values of the state variables at time t, $x(t)$.

(2) the values of a set of informational variables, say $v_j(t)$ for $j = 1, \ldots, m$, which, in general, vary with time (these may depend upon or include driving variables).

(3) the values of a set of parameters or constants, say p_k, $k = 1, \ldots, s$, which do not vary with time.

(4) and time itself.

The change in the ith state variable $\Delta x_i(t)$ at time t over the time interval Δt may be functionally written as

$$\Delta x_i(t) = F_i[\underset{\sim}{x}(t), \underset{\sim}{v}(t), \underset{\sim}{p}, t, \Delta t] \cdot \Delta t,$$

where the function F_i is the change per unit time in the state variable x_i. Note that if the system modeled is to be represented by *differential* as opposed to *difference* equations, then the dependence of F_i upon Δt should not exist.

Given the initial values of the state variables at time $t = t_0$, that is $\underset{\sim}{x}(t_0)$, and the changes in the state variables $\Delta \underset{\sim}{x}(t)$, we can find the state of the system at any time $t_m = t_0 + m\Delta t$ for $m = 0, 1, 2, \ldots, M$. The state of the system at any time t_M is iteratively computed as

$$\underset{\sim}{x}(t_M) = \underset{\sim}{x}(t_0) + \Delta t \cdot \sum_{m=0}^{M-1} F_i[\underset{\sim}{x}(t_m), \underset{\sim}{v}(t_m), \underset{\sim}{p}, t_m, \Delta t]$$

In order to simulate biological systems, we postulate the following three principles.

(1) A biological system can be viewed as a collec-

tion of smaller subsystems. (Indeed some

systems might consist of a single subsystem.)

(2) A change of state in any subsystem must result

from the flow of material or energy between

compartments contained in that subsystem.

(3) The identity of the material or energy flowing

in any subsystem must maintain its physical

identity throughout the subsystem.

As a result of the second postulate, we have further

required that the change of state of any particular com-

partment be expressed as the algebraic sum of the flows

to or from that compartment. Let the *net* flow per unit

time from compartment i to compartment j be represented by

$$f_{ij} = f_{ij}[\underset{\sim}{x}(t), \underset{\sim}{v}(t), \underset{\sim}{p}, t].$$

Note that $f_{ij} = -f_{ji}$, that is the net flow into compart-

ment j, is reflected by a corresponding loss from

compartment i, and by necessity the identity of the

material flowing must remain unique. Therefore, expanding

upon our formulation of the solution of the initial

value problem, we find that the rate of change of

material in some compartment i, above expressed as

$F_i[\underset{\sim}{x}(t), \underset{\sim}{v}(t), \underset{\sim}{p}, t]$, is the sum of the net flows from

each of the associated compartments. Formally this

requires that $F_i = \underset{j \varepsilon S}{\Sigma} f_{ij}$, where S is the set of

compartments which are coupled to compartment i by

flows.

A nine-compartment system comprised of two sub-
systems is illustrated in Fig. 1.1-1. The compartments
are represented by the boxes. Material or energy flows
are represented by solid arrows between the compartments.
A flow which is always in one direction is represented
by a single-headed arrow, such as the flow from compart-
ment 1 to compartment 3. Flows in which the net flow
may be in either direction are represented by a double-
headed arrow, such as the flow between compartment
2 and compartment 3. Note that there are no material
flows between compartments in separate subsystems.
Informational flows are represented by dotted arrows.
The rate of flow between compartment 5 and compartment
7 for example, is controlled by the amount of material
in compartment 3. These informational flows are
represented by $\underset{\sim}{v}(t)$ in our mathematical formulation.
We may further identify compartment 1 as a source,
provided the flow from 1 to 3 does not depend upon the
quantity of material in 1, and likewise identify
compartment 8 as a sink.

Fig. 1.1-1. A nine-compartment system comprised of two subsystems illustrating actual flows (solid arrows) and informational flows (dotted arrows).

We have now identified the elements necessary to specify a computer simulation of a compartmental-flow model. We require,

(1) initial values for the state variables.

(2) mathematical expressions which calculate the net flows between compartments.

(3) mathematical formulas which calculate the informational flows.

(4) the identification and values for parameters and constants used in (2) and (3).

(5) the starting and final times over which the simulation is to be run and the time step for the numerical solution of the initial value problem.

SIMCOMP is designed so that these five items are easily specified by the user. SIMCOMP organizes the flow expressions into difference equations and provides the solution. SIMCOMP further requires the user to specify what information is to be printed and plotted. The syntactical definitions for writing SIMCOMP programs are explained in section 2. The mathematical formulation of any compartmental-flow simulation, by first specifying the above five items, perhaps with the aid of a flow diagram, should precede the formulation of a SIMCOMP program.

1.2 Event Simulations.

While many ecological processes can be visualized as material or energy transfers between compartments, some processes are not so easily represented. Processes in the following categories can often be most easily described by an event-oriented simulation:

- Processes involving discrete-valued variables.

- Processes which can be visualized as a queuing problem.

- Processes which do not occur uniformly through time.

- Stochastic processes.

An event in SIMCOMP can be formally defined by specifying the following two items:

(1) A computation or set of computations, referred to as the *action* of the event, which represents the effect of the event on the system.

(2) A specification of the *time of occurrence* of the event.

All events will be assumed to have the above two attributes. The development of an event-oriented simulation model will then involve abstracting from real-world phenomena, processes which can be completely described by specifying the action of the process on the system and the timing of the process. For example, if a birth process is to be simulated, the action of the

process would be to increase the number of individuals in the population. The time of occurrence of the birth would also have to be determined.

There exists two methods by which the time of occurrence of an event can be determined: events generated within the model and events fed to the model from the outside world. The first method of event scheduling is termed *internal* or *endogenous* while the second method is termed *external* or *exogenous*. The difference between the types is that endogenous events are triggered by the explicit reaction of the model to its operations, i.e., the model generates internal events as it progresses, while exogenous events are fed to the model from an external data source.

Changes which take place in the state of the system when an event occurs are termed *actions*. Central to the concept of an event is that an action requires zero-simulated time to occur. This is the crucial difference between discrete-event and continuous-time simulations. In discrete-event simulations, state-changes take place only at specified points in simulated time at which interactions between system components occur. In continuous-time simulations, interactions and state-changes take place continuously. To model continuous changes, numerical integration procedures must be employed, but are not required for discrete-event simulations.

2. SIMCOMP PROGRAMMING.

SIMCOMP is a FORTRAN-like language designed to implement both continuous and event simulations as described in section 1. Continuous and discrete-event simulations may be combined in the same simulation.

The fundamental elements of a continuous-variable simulation are *flow definitions*. Likewise *event definitions* are the fundamental elements of discrete simulations. Storage allocation for globally defined variables (i.e., variables accessible by all portions of the simulation) is provided by means of *storage declarations*. FORTRAN subroutines and functions may be supplied by the user. The *source section* of a SIMCOMP simulation is specified by the inclusion of any or all of the above statement types. The format and usage of SIMCOMP source statements are described in section 2.1.

A SIMCOMP simulation source section is a mixture of FORTRAN statements and SIMCOMP processor directives. This manual assumes a basic knowledge of FORTRAN programming and the user is referred to any good instructional FORTRAN manual such as "Computer Programming - FORTRAN IV" (Anderson 1966). The SIMCOMP processor produces code which is compiled by Control Data Corporation's FORTRAN Extended Version 3.0 compiler. It is recommended that all FORTRAN coding contained in a SIMCOMP

source program conform to the specifications in the
"FORTRAN Extended Reference Manual" (Control Data
Corporation 1973).

The initial values of state variables and parameters
are specified in the SIMCOMP *data section*. Requests for
tabular and graphic output are also included in the data
section. The format and usage of SIMCOMP data section
specifications are described in section 2.2. The data
section is read in and processed by the SIMCOMP-generated
simulation program. The sequence of operations in the
processing of the source and data sections is outlined
in Fig. 2-1.

Fig. 2-1. Sequence of operations in a typical SIMCOMP job. Dotted blocks
indicate user-supplied portions.

The detection and notification to the user of execution errors is a task usually assigned to the computer operating system. SIMCOMP provides an execution time debugging facility which recovers control from the operating system when an error is detected. A mnemonic dump of user variables along with an explanation of the nature of the error is printed in the output. Section 2.3 describes how this information may be utilized in debugging programs.

2.1 <u>Source</u> <u>Program</u>.

SIMCOMP source programs are a mixture of _SIMCOMP processor directives_ and _FORTRAN statements_. The SIMCOMP compiler is actually a pre-processor which converts SIMCOMP source language statements into segments of FORTRAN compilable code.

SIMCOMP recognizable processor directives are generally comprised of a key word followed by a period. These statements may be contained anywhere in columns 1 through 72. FORTRAN statements included in the text must begin in or after column 7. Columns 1 through 5 are used for statement labels and columns 73 through 80 are ignored. Column 6 is used for statement continuation. _FORTRAN statement labels may take on any value with the exception of five digit labels beginning with 9._

SIMCOMP reserves certain variable names as attributes of the system. Any of these variables may be used (or altered at the user's discretion) in the computation at any time. These variables and their meaning are shown in Table 2.1-1.

Table 2.1-1.   Reserved variable names (simulation control variables).

Variable	Meaning
X(i) where $1 \leq i \leq 999$	Current amount of material in compartment i.
TIME	Current simulated time.
TSTRT	Starting time of the simulation.
TEND	Ending time of the simulation.
DT	Integration step size.
DTPR	Time step between print-outs.
DTPL	Time step between plotted values.
DTFL	Time step between flow print-outs.
FLOW	Value of the currently computed flow.

The user is cautioned against the use of any variable name beginning with the letter X. Variables which are internal to the operation of the SIMCOMP system use the convention of beginning with the letter X. This avoids potential conflicts between the user-supplied code and the system routines. This precaution deserves special cognizance for "canned" FORTRAN subroutines where such variables might be used.

2.1.1 Parameter declarations.

Parameter declarations are used for two purposes. These are (i) storage allocation and (ii) stochastic function definition. All parameter declarations consist of a key word followed by a period, followed by a list of names delimited by commas of the following form:

$$\text{key word.} \quad name_1, name_2, \ldots, name_n$$

The key word may begin in any column. The entire statement should be contained in or before column 72. Parameter declaration statements may not be continued on successive cards. As many parameter declaration statements may be included as are required. Parameter declaration statements can appear anywhere in the source program with the following exceptions:

(1) within the text of a flow.

(2) within an event routine or subprogram.

Variable storage allocation.

Storage allocation statements consist of statements of the following form:

$$\text{STORAGE.} \quad var_1, var_2, \ldots, var_n$$
$$\text{INTEGER.} \quad var_1, var_2, \ldots, var_n$$
$$\text{REAL.} \quad var_1, var_2, \ldots, var_n$$

The names of variables in the variable declaration list may be from 1 to 5 alphanumeric characters in length and must begin with a letter other than X.

Variables which fall in the following categories should be declared in a variable storage allocation statement.

(1) Variables which are subscripted.

(2) Variables whose values are to be assigned via data assignment statements in the data section (refer to section 2.2.1).

(3) Variables where values are to be printed or plotted via PRINT.[1] or PLOT. requests. (refer to section 2.2.2).

(4) Variables whose values are computed in flows and are used in events or subprograms and vice versa.

(5) Variables whose implicit type must be altered, i.e., from integer to real or from real to integer.

STORAGE., INTEGER., and REAL. statements can be thought of as FORTRAN COMMON, INTEGER, and REAL declaration statements. This is true with the exception that any variable declared in an INTEGER. or REAL. statement is treated as though the variable were also declared in a STORAGE. statement. As such any variable declared in one of the three storage allocation statements can be considered to be globally defined in all segments of the simulation. *All events and subprograms, in addition*

---

[1] Note that the period is part of the command verb.

to all flows, have access via its mnemonic name to the value of any declared variable. The maximum number of dimensions allowed for any subscripted variable is three.

Example 2.1.1-1. Storage allocation statements.

```
STORAGE. A,B,P(3),Q(2,3),INDEX
REAL. M(3),N,L(2,2)
INTEGER. D,E(2,3,2),F,B,P(6),L
```

If a variable is declared more than once, the last declared mode of the variable is assumed. In the above example variable "B" would be assumed type integer. If a variable is dimensioned more than once, the last declared dimensions hold. In the above example the variable "P" is assumed to be an integer one-dimensional array with six locations. Once a variable is dimensioned, the dimensionality of the variable remains in force regardless of changes in type. The variable "L" above is assumed to be an integer two-dimensional array with four locations (two by two).

Primary and secondary class storage.

Normally, the initial values of all user-declared variables are printed in the output after the data section has been processed, prior to the start of the simulation. User-declared variables can be segregated into two classes of variables by the following convention.

Any user-declared variables named in storage allocation statements in the normal manner will be considered a primary-class variable. All primary-class variables will be printed in the initial-conditions output unless otherwise requested. Secondary-class variables are prefaced by an asterisk. Normally, secondary-class variables will not be printed in the initial-conditions output. Secondary-class variables are treated just as primary-class variables in all other respects. This feature is useful especially in the case of large arrays whose initial conditions are not of interest, thus minimizing the amount of output produced in the initial-conditions output. Data section commands which will alter the normal procedure taken for selecting variables for printing in the initial-conditions output are described on page 2.2.2-17. An example of secondary-class variables is also presented in section 2.2.2.

Stochastic function definitions.

Stochastic function definition statements consist of statements of the following form:

UNIFORM. $name_1$, $name_2$, ... , $name_n$

NORMAL. $name_1$, $name_2$, ... , $name_n$

EXPONENT. $name_1$, $name_2$, ... , $name_n$

LOGNORMAL. $name_1$, $name_1$, ... , $name_n$

The names contained in the variable list must contain from 1 to 5 alphanumeric characters beginning with a

letter. Variable names starting with the letter X should not be used. Similarly variable names which are implicitly type integer should not be used. Continuation cards are not allowed. As many stochastic function definition statements as required may be included.

Each entry in the list does not actually allocate storage to the named variable, but generates a function subprogram of that name. The function is called by using the variable name in an arithmetic expression. In the expression the variable name must be followed by an argument list containing the correct number of parameters which specify the particular distribution function. The parameters in the argument lists must be real-valued constants or variables. Each call returns a value from the indicated distribution as the value of the function. The number of parameters and their meanings for each of the distributions are given in Table 2.1-2.

Table 2.1-2. Stochastic variable parameters.

Distribution Function	No. of Parameters	Meaning of Parameters
Uniform	2	(1) Minimum value. (2) Maximum value.
Normal	2	(1) Mean value. (2) Standard deviation.
Exponential	1	(1) Expected value.
Lognormal	2	(1) Mean value. (2) Standard deviation.

Example 2.1.1-2.   A flow simulation containing stochastic parameters.

```
NORMAL. V
STORAGE. VMEAN,VSTD,P
(1-2). RV=V(VMEAN,VSTD)
       FLOW=RV*(P-X(2))
78g     end-of-record separator 2/
VMEAN=0.01 $ VSTD=0.01 $ P=20. $ X(1)=100. $ X(2)=5. $
TSTRT=0. $ TEND=100. $ DT=1. $
PLOT. (X(1)=1),(X(2)=2)
```


The parameter RV in the above example might represent the value of some variable which was experimentally determined to have a mean value of 0.01 with a standard deviation of 0.01.   The above flow would be computed

---

2/ An end-of-run separator is a single card with a 7-8-9 multipunched in column 1.

using a randomly sampled value from a normal distribu-

tion with the given mean and standard deviation at each

time step of the simulation. Refer to section 2.1.2

for a description of flow definitions.

A description of each of the above distributions

and the method used for their generation is given by

Naylor et al. (1966).

Example 2.1.1-3. An event simulation using stochastic parameters.

```
STORAGE. RMIN,RMAX,NO,TEXP,FINAL
UNIFORM. SIZE
EXPONENT. TDELT
        EVENT MIGRT
        TNEXT=TIME+TDELT(TEXP)
        IF(TNEXT.GT.FINAL) RETURN
        NO=NO+SIZE(RMIN,RMAX)
        CALL EVENT(5HMIGRT,TNEXT,1)
        RETURN
        END
```

7_{89} *end-of-record separator*

```
RMIN=10. $ RMAX=35. $ NO=0 $ TEXP=1. $ FINAL=255. $
TSTRT=0. $ TEND=365. $
EVENT. MIGRT,245.,1
PLOT. (NO=N)
```

PLOT NO. 1

TIME

This example might simulate the immigration of a species of animal during the time interval from day 245 to day 255. The time interval between arrivals of groups of animals was assumed to be exponentially distributed with an expected value of one. The number of animals per group was assumed to be a uniformly distributed random variable in the range from 10 through 35. The total number of animals which have arrived is contained in the variable NO. Refer to section 2.1.3 for a description of event definitions.

2.1.2 Flow definitions

Flows or material transfers between compartments, named $X(j)$ where $1 \leq j \leq 999$, in a subsystem are computationally defined in flow definitions. A flow definition is comprised of a flow definition label followed by a series of one or more FORTRAN statements which compute the flow rate. The reserved variable FLOW should be set to the computed value of the flow rate. Flow definitions in general take the following form:

(phrase - phrase).

executable FORTRAN statements

In the above, the terms "phrase" are each one of the following forms:

(1) n where n is an integer constant.

(2) $v = n_1, n_2$ where v is a simple integer variable and n_1 and n_2 are integer constants.

(3) $v = n_1, n_2, n_3$ where v is a simple integer variable and the n_i are integer constants, $i = 1, 2, 3$.

(4) $v_1 = n_1 * v_2 \pm n_2$ where the v_i are simple integer variables and the n_i are integer constants, $i = 1, 2$.

The above phrases specify the indices of the source and destination state variable compartments between which a flow occurs. The system allows for a maximum

of 999 compartments (i.e., X(1) through X(999)). The
maximum number of flows which may be defined is 9999
or is limited by the amount of central memory core
storage available. A flow definition label may begin
in any column and must be completed in or before column
72. Any nonblank characters following the period in
or before column 72 are assumed to be an executable
FORTRAN statement.

Constant phrases.

If either of the phrases in a flow definition
label are of the form (1) above, n must be an integer
constant in the range $1 \leq n \leq 999$. The following flow
label would define a flow from compartment X(3) to
compartment X(239).

(3-239).

executable FORTRAN statements

Iterative phrases.

Phrases of the form (2) and (3) define flows
iteratively. If the mathematical form of a series of
flows is identical, perhaps differing only in the
values of parameters used in the computation, the
flows can be economically written. The phrases of

forms (2) and (3) correspond in operation to the iteration phrase of a FORTRAN DO-loop. Phrases of the form (2) would indicate a series of compartments n_1, $n_1 + 1$, $n_1 + 2$, ... , n_2. These are the values that the integer variable v takes on. Admissable values of the constants n_1 and n_2 must satisfy $1 \leq n_1 \leq n_2 \leq 999$. Phrases of the form (3) would indicate a series of compartments n_1, $n_1 + n_3$, $n_1 + 2 * n_3$, $n_1 + 3 * n_3$, ... , $n_1 + m * n_3$ where m is the smallest value such that $n_1 + m * n_3 \geq n_2$. Admissable values of the constants n_1, n_2 and n_3 must satisify $1 \leq n_1 \leq n_2 \leq 999$ and $n_1 + m * n_3 \leq 999$. Source and destination compartment phrases may contain any combination of forms (1), (2), and (3). The integer valued variable v must be a simple integer variable containing from 1 to 5 characters. The following flow declarations illustrate some of the possible combinations. A flow diagram of the flows defined by each declaration is included with each case.

Case 1.

(733 - I = 83, 85).

executable FORTRAN statements

Case 2.

(KK = 1, 5, 2 - 20).

executable FORTRAN statements

Case 3.

(IFROM = 1, 7, 3 - ITO = 998, 999).

executable FORTRAN statements

Computational phrases.

Phrases of the form (4) must be used in conjunction with iteration phrases of the forms (2) and (3). The variables v_1 and v_2 must be simple integer variables containing five or fewer alphanumeric characters. The variable v_2 must be the same variable used in

the other half of the flow definition. The constants

n_1 and n_2 must be simple integer constants. If either

of the constants n_1 and n_2 are chosen to have the

value zero, the zero must be written. That is, a

computational phrase must appear exactly as specified

in form (4). The values of the constants n_1 and n_2

must be chosen such that the values of v_1 satisfy

$1 \leq v_1 \leq 999$. The following declarations illustrate

the usage of computational phrases of form (4). A

flow diagram of the flows defined follows each case.

Case 1.

$(I = 1, 3 - J = 2 * I + 2)$.

executable FORTRAN statements

Case 2.

```
┌──────┐          ┌──────┐
│ 100  │─────────▶│ 101  │
└──────┘          └──────┘

┌──────┐          ┌──────┐
│ 102  │─────────▶│ 103  │
└──────┘          └──────┘

┌──────┐          ┌──────┐
│ 104  │─────────▶│ 105  │
└──────┘          └──────┘
```

$(M = 1 * N - 1 - N = 101, 105, 2)$.

executable FORTRAN statements

Case 3.

```
┌──────┐          ┌──────┐
│ 100  │─────────▶│ 200  │
└──────┘          └──────┘

┌──────┐          ┌──────┐
│ 101  │─────────▶│ 202  │
└──────┘          └──────┘
```

$(I1 = 100, 101 - I2 = 2 * I1 + 0)$.

executable FORTRAN statements

The statements which follow a flow definition label can be any executable FORTRAN statements with the following restrictions.

(1) FORTRAN statement labels containing five numeric characters beginning with "9" should not be used. Any FORTRAN statement label which is not of the form 9DDDD, where the D's are any digits, may be used.

(2) FORTRAN transfer of control statements (i.e.,
conditional or unconditional jumps) should
not transfer control to statements not contained
within the range of the current flow definition
label. The range of a flow definition label is
defined as all executable FORTRAN statements
following the flow definition label prior to
encountering (i) another flow definition label,
(ii) a parameter declaration statement (refer
to section 2.1.1), (iii) a SUBROUTINE, FUNCTION,
or EVENT statement (refer to sections 2.1.3
and 2.1.4), or (iv) the end of the source program.

(3) The reserved system variable FLOW is set
equal to the computed value for the flow rate.
The value of FLOW does not have to be set via
an arithmetic replacement statement. FLOW
may be passed as a formed parameter to a
subprogram where its value is set. If within
a flow definition FLOW is not assigned a
value, its value is flagged as INDEFINITE
and a fatal error will occur (refer to section
2.3).

Construction of flow simulations.

The design and construction of flow-oriented
continuous variable simulations might be described by
the following steps.

(1) Construct a flow diagram of the system to be simulated.

(2) Develop mathematical equations which will compute the values for each of the flow rates.

(3) Program the flow definition statements.

(4) Execute, debug, and evaluate the output of the simulation.

The following soil water model, taken from Smith (1971), is presented as an example of a compartmental flow model. The model is designed to simulate the following processes:

(1) infiltration of surface water.

(2) surface water runoff.

(3) transfer of soil water from unsaturated to saturated storage elements.

(4) soil water drainage.

Evapotranspiration is not considered in the model. This model is presented primarily to illustrate the implementation of the model in SIMCOMP. A complete discussion of the theory and performance of the model is presented in Smith (1971). A flow chart of the model is presented in Fig. 2.1.2-1. The following verbal description of the operation of the model is excerpted from the above mentioned report.

Fig. 2.1.2-1. Flow chart of the volumetric threshold infiltration model.

The compartments X(1), X(2), and X(3) represent the volumes for depression storage, unsaturated soil water storage, and the total volume between the unsaturated storage and saturation. The input to the system is the rain rate R1. The outputs of the system are the runoff rate Y1 and the drainage rate Y3. Ordinarily, evapotranspiration would draw water from compartment X(2). The flow between X(1) and X(2) is the actual infiltration rate Y12 and the flow between X(2) and X(3) is also the actual infiltration rate Y23.

If the rain rate is greater than the potential infiltration rate FP or if there is ponded water at the surface, i.e., X(1) is greater than zero, Y12 is equal to the potential infiltration rate. If the rain rate is nonzero and is less than the potential infiltration rate and X(1) is zero, Y12 is equal to the rain rate.

If X(1) is less than the surface storage capacity K1, the runoff rate Y1 is zero. When X(1) attempts to exceed K1, the runoff rate is equal to the difference between the rain rate and the infiltration rate.

When X(2) is less than the capacity of unsaturated storage K2, Y23 is zero. When X(2) approaches K2, steady state is reached for that storage, i.e., input equals output, and Y23 is equal to Y12. When X(3) is greater than zero, the drainage rate Y3 is equal to the saturated hydraulic conductivity (the final infiltration rate, FC).

A listing of the source and data sections of the model and the output produced during a 9-hour simulation comprised of two rainfall events is presented in the following example. The graphs presented were reproduced from microfilm which was generated by SIMCOMP.

Example 2.1.2-1. A sample simulation illustrating flows.

```
STORAGE. RAIN(2.24)
STORAGE. A.FC
REAL. N.K1.K2.K3
STORAGE. R1.Y1.Y12.Y23.Y3.FP
C....COMPARTMENT DEFINITIONS.
C
C     X(1)          DEPRESSION STORAGE.
C     X(2)          UNSATURATED STORAGE.
C     X(3)          SATURATED STORAGE.
C     X(10)         TOTAL RUNOFF.
C     X(20)         DEEP STORAGE.
C     X(999)        SOURCE OF RAINFALL.
C
C....VARIABLE DEFINITIONS.
C
C     K1            VOLUME OF DEPRESSION STORAGE.
C     K2            VOLUME OF UNSATURATED STORAGE.
C     K3            VOLUME OF SATURATED STORAGE.
C     FC            SATURATED HYDRALIC CONDUCTIVITY
C                      (FINAL INFILTRATION RATE).
C     N             POTENTIAL INFILTRATION EXPONENT.
C     A             POTENTIAL INFILTRATION COEFFICIENT.
C     R1            RAIN RATE.
C     FP            POTENTIAL INFILTRATION RATE.
C  )  Y1            RUNOFF RATE.
C     Y12           ACTUAL INFILTRATION RATE (X(1) TO X(2)).
C     Y23           ACTUAL INFILTRATION RATE (X(2) TO X(3)).
C     Y3            DRAINAGE RATE.
C     RAIN          RAIN RATE DATA RECORD.
C
C...THE RAIN RATE IS LINEARLY INTERPOLATED FROM DATA.
(999-1). R1=ALINT2(TIME.IFCK.RAIN)
     FLOW=R1
C...INFILTRATION TO UNSATURATED STORAGE.
(1-2). FP=A*(K2+K3-X(2)-X(3))**N+FC
     Y12=FP
     IF (X(1).LE.0.) Y12=AMIN1(R1.FP)
     IF (Y12*DT.GT.X(1)) Y12=X(1)/DT
     FLOW=Y12
C...RUNOFF WHEN THE CAPACITY OF DEPRESSION STORAGE IS EXCEEDED.
(1-10). Y1=0.
     IF (X(1).GT.K1) Y1=AMAX1(R1-Y12.0.)
     IF ((Y12+Y1)*DT.GT.X(1)) Y1=(X(1)-Y12*DT)/DT
     FLOW=Y1
C...INFILTRATION TO SATURATED STORAGE.
(2-3). Y23=0.
     IF (X(2).GT.K2) Y23=Y12
     IF (Y23*DT.GT.X(2)) Y23=X(2)/DT
     FLOW=Y23
C...DRAINAGE WHEN THE SOIL IS SATURATED.
(3-20). Y3=AMIN1(Y23.FC)
     IF (X(3).GT.0.) Y3=FC
     IF (Y3*DT.GT.X(3)) Y3=X(3)/DT
     FLOW=Y3
```

$^7_8{}_9$ end-of record separator

```
K1=0.1 $ K2=1.60 $ K3=0.15 $ A=0.65 $ N=1.19 $ FC=0.83 $
X=3*0. $ X(10)=0. $ X(20)=0. $ X(999)=1000. $
TSTRT=0. $ TEND=9. $ DT=0.02 $ DTPR=0.1 $
R1=0. $ Y1=0. $ Y3=0. $ Y12=0. $ Y23=0. $ FP=0. $
RAIN=0.,0.,0.1,2.6,0.2,5.6,0.3,6.0,0.4,5.8,0.5,4.6,0.6,3.4,0.7,2.0,0.8,1.0,
     0.9,0.4,1.0,0.,5.,0.,5.05,1.3,5.1,2.8,5.15,3.0,5.2,2.9,5.25,2.3,5.3,1.7,
     5.35,1.0,5.4,0.5,5.45,0.2,5.5,0. $
PRINT.
PRINT. R1,FP,Y1,Y12,Y23,Y3
TITLE.        RAINFALL RATE GENERATED FROM DATA.
PLOT. (R1)
TITLE.   DEPRESSION STORAGE = 1,  RUNOFF RATE = Y,  ACTUAL INFILTRATION RATE = A
PLOT. (Y1=Y),(X(1)=1),(Y12=A)
TITLE.   UNSATURATED = 2,  SATURATED = 3,  POTENTIAL INFILTRATION RATE = P
PLOT. (X(2)=2),(X(3)=3),(FP=P)
TITLE.   DEPRESSION STORAGE = 1,  RUNOFF RATE = Y,  ACTUAL INFILTRATION RATE = A
PLOT. (Y1=Y),(X(1)=1),(Y12=A)[0.,2.]
TITLE.   DEPRESSION STORAGE = 1,  RUNOFF RATE = Y,  ACTUAL INFILTRATION RATE = A
PLOT. (Y1=Y),(X(1)=1),(Y12=A)[4.5,6.]
TITLE.   UNSATURATED = 2,  SATURATED = 3,  POTENTIAL INFILTRATION RATE = P
PLOT. (X(2)=2),(X(3)=3),(FP=P)[0.,2.]
TITLE.   UNSATURATED = 2,  SATURATED = 3,  POTENTIAL INFILTRATION RATE = P
PLOT. (X(2)=2),(X(3)=3),(FP=P)[4.5,6.]
FILM.
```
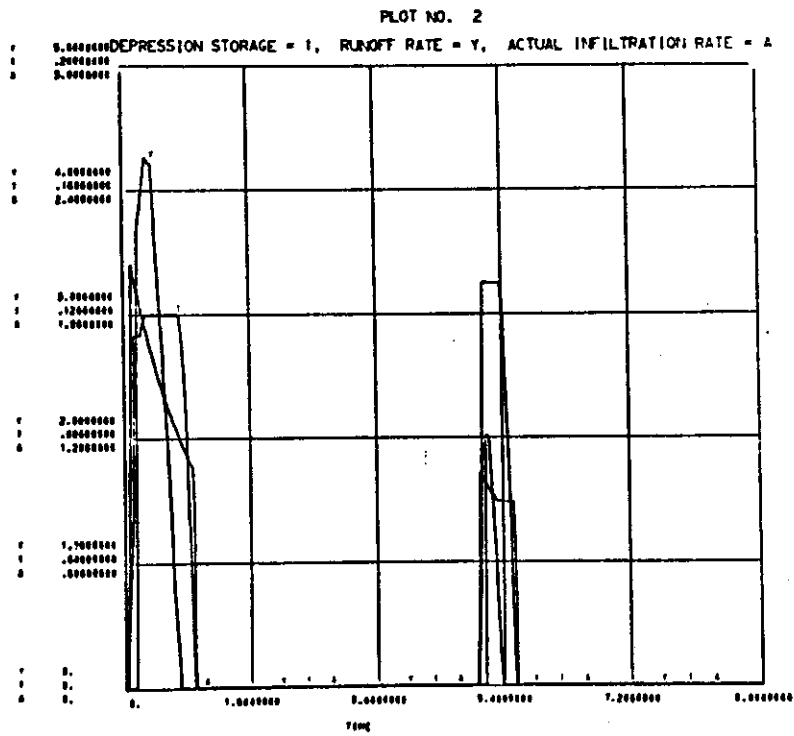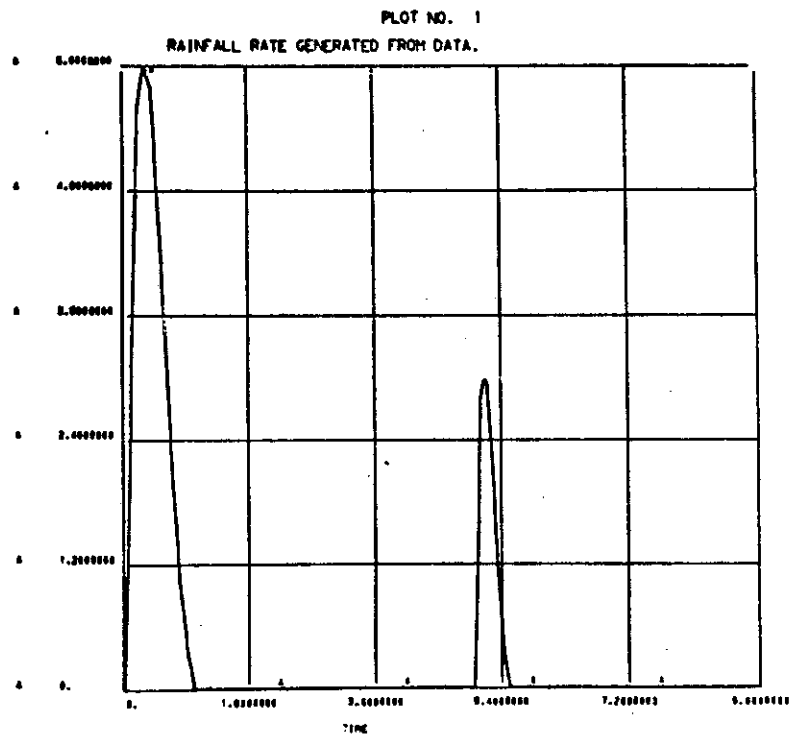
SIMULATION RESULTS

(partial listing)

```
TIME = 0.
       X(999) = 1000.00000      X(1) =           0.      X(2) =           0.      X(10) =           0.
       X(3) =           0.      X(20) =           0.      R1 =           0.      FP =           0.
       Y1 =           0.        Y12 =           0.        Y23 =           0.      Y3 =           0.

TIME = .100000000
       X(999) = 999.896000      X(1) = .416000000E-01   X(2) = .624000000E-01   X(10) =           0.
       X(3) =           0.      X(20) =           0.      R1 = 2.00000000      FP = 2.06831609
       Y1 =           0.        Y12 = 1.56000000        Y23 =           0.      Y3 =           0.

TIME = .200000000
       X(999) = 999.404000      X(1) = .113176426      X(2) = .297248000      X(10) = .185575574
       X(3) =           0.      X(20) =           0.      R1 = 5.60000000      FP = 1.87494158
       Y1 = 3.72508042        Y12 = 1.87494158        Y23 =           0.      Y3 =           0.

TIME = .300000000
       X(999) = 998.820000      X(1) = .120000000      X(2) = .475640767      X(10) = .584359233
       X(3) =           0.      X(20) =           0.      R1 = 6.00000000      FP = 1.72540424
       Y1 = 4.19459576        Y12 = 1.72540424        Y23 =           0.      Y3 =           0.

TIME = .400000000
       X(999) = 998.232000      X(1) = .120000000      X(2) = .640008860      X(10) = 1.00799114
       X(3) =           0.      X(20) =           0.      R1 = 5.80000000      FP = 1.59111732
       Y1 = 4.20888268        Y12 = 1.59111732        Y23 =           0.      Y3 =           0.

TIME = .500000000
       X(999) = 997.724000      X(1) = .120000000      X(2) = .791777277      X(10) = 1.36422272
       X(3) =           0.      X(20) =           0.      R1 = 4.60000000      FP = 1.47044433
       Y1 = 3.12955567        Y12 = 1.47044433        Y23 =           0.      Y3 =           0.

TIME = .600000000
       X(999) = 997.336000      X(1) = .120000000      X(2) = .932220648      X(10) = 1.61177935
       X(3) =           0.      X(20) =           0.      R1 = 3.40000000      FP = 1.36196648
       Y1 = 2.03803352        Y12 = 1.36196648        Y23 =           0.      Y3 =           0.

TIME = .700000000
       X(999) = 997.080000      X(1) = .120000000      X(2) = 1.06240353      X(10) = 1.73751647
       X(3) =           0.      X(20) =           0.      R1 = 2.00000000      FP = 1.26445921
       Y1 = .735640793        Y12 = 1.26445921        Y23 =           0.      Y3 =           0.

TIME = .800000000
       X(999) = 996.940000      X(1) = .116462491      X(2) = 1.18359000      X(10) = 1.75993871
       X(3) =           0.      X(20) =           0.      R1 = 1.00000000      FP = 1.17687545
       Y1 =           0.        Y12 = 1.17687545        Y23 =           0.      Y3 =           0.

TIME = .900000000
       X(999) = 996.876000      X(1) = .475565142E-01   X(2) = 1.29650478      X(10) = 1.75993871
       X(3) =           0.      X(20) =           0.      R1 = .400000000      FP = 1.09833736
       Y1 =           0.        Y12 = 1.09833736        Y23 =           0.      Y3 =           0.

TIME = 1.00000000
       X(999) = 996.860000      X(1) = .170002901E-14   X(2) = 1.38006129      X(10) = 1.75993871
       X(3) =           0.      X(20) =           0.      R1 = .052651203E-13   FP = 1.03000669
       Y1 =           0.        Y12 = .800000000E-01     Y23 =           0.      Y3 =           0.

TIME = 1.10000000
       X(999) = 996.860000      X(1) =           0.      X(2) = 1.38006129      X(10) = 1.75993871
       X(3) =           0.      X(20) =           0.      R1 =           0.      FP = 1.02906174
       Y1 =           0.        Y12 =           0.        Y23 =           0.      Y3 =           0.
```

GRAPHICAL SIMULATION RESULTS   07/23/73   19.11.52.

GRAPH NO.	GROUP	GROUP RANGE DECLARATION	DEPENDENT VARIABLE(S)	PLOTTED CHARACTER	INDEPENDENT VARIABLE	INDEPENDENT VARIABLE RANGE DECLARATION
1	1		NJN NJF NJ	N F J	YEAR	
2	1		NAN NAF NA	N F A	YEAR	
3	1		NJ NA N	J A N	YEAR	
4	1		NJ	J	YEAR	
5	1		NA	A	YEAR	
6	1		N	N	YEAR	

PLOT NO. 1

RAINFALL RATE GENERATED FROM DATA.

PLOT NO. 2

DEPRESSION STORAGE = 1,   RUNOFF RATE = Y,   ACTUAL INFILTRATION RATE = A

PLOT NO. 3

UNSATURATED = 2,   SATURATED = 3.   POTENTIAL INFILTRATION RATE = P

PLOT NO. 4

DEPRESSION STORAGE = 1,   RUNOFF RATE = Y,   ACTUAL INFILTRATION RATE = A

PLOT NO. 5

DEPRESSION STORAGE = 1, RUNOFF RATE = Y, ACTUAL INFILTRATION RATE = A

TIME

PLOT NO. 6

UNSATURATED = 2, SATURATED = 3, POTENTIAL INFILTRATION RATE = P

TIME

PLOT NO. 7

UNSATURATED = 2, SATURATED = 3, POTENTIAL INFILTRATION RATE = P

TIME

2.1.3 Event definitions.

An event in the SIMCOMP language is defined as a set of computations which may be scheduled for execution at any instant during simulated time. Event routines are essentially FORTRAN subroutines which are called by the executive routine at requested times. The format of an event routine is:

```
EVENT name
       .
       .
       .
FORTRAN statements
       .
       .
       .
END
```

All statements must begin in or after column 7. Columns 72 through 80 are ignored. Column 6 may be used for statement continuation. The name of the event *name* must contain from one to five alphanumeric characters starting with a letter. Event names beginning with the character X should be avoided. The FORTRAN statements which represent the computations in the event must be followed by an END card. All variables which have been declared in parameter declaration statements (refer to section 2.1.1), in addition to the system reserved variables, may be considered present in the event and may be used in the computations. A maximum of 100 different events can be *defined* in a simulation.

Event scheduling.

SIMCOMP simulations are executed under the control of an executive routine. This executive routine has the responsibility of stepping the simulation through time. In addition to scheduling and passing control to a number of system-defined events such as printing output, saving values of variables for plotting, and updating the state variables if flows are included, the SIMCOMP executive keeps a dynamic list of all user-defined events scheduled to occur and their time of occurrence. This list is termed the event stack.

An event can be scheduled to occur either exogenously (externally) or endogenously (internally). Exogenous events are defined as those events which are scheduled prior to the start of simulated time. By including an exogenous event request card in the data (refer to section 2.2.3), an event, its time of occurrence, and a priority is entered into the event stack. Exogenous event request cards have the following format.

EVENT. name, time, priority

Endogenous events are defined as those events which are scheduled dynamically during the course of a simulation. An event is placed in the event stack by a FORTRAN call to the system event schedule in the following format,

CALL EVENT(mHname,time,priority)

where,

m	is a character count of the number of characters contained in the name of the event $(1 \leq m \leq 5)$.
name	is the name of the event routine (left justified). The term mHname is a FORTRAN hollarith constant.
time	is a real-valued variable or constant containing the value of simulated time at which the event is to occur.
priority	is an integer variable or constant in the range 1 through 512.

After a call of the above form is made and the current simulated time, TIME, becomes equal to the scheduled time of occurrence of the event, the event is called and executed. When an event is called by the executive routine, the corresponding entry in the event stack is purged. The priority of an event is used as a tie breaker if more than one event is scheduled to occur at the same time. A priority of 1 is highest (first to occur) and 512 is lowest (last to occur). If the value of a priority is outside the range 1 through 512, a priority of 512 is assumed. If two or more events of the same priority are scheduled to occur at the same time, the first to have been scheduled is the first to occur. Additionally, if the same event is scheduled

to occur more than once at identical times, the second
and subsequent requests are ignored. The maximum number
of events that can be scheduled at any one time is
limited by the amount of core available to the job. If
an attempt to schedule a nonexistent event is made, a
diagnostic is issued and the simulation is terminated.

Once an event has been put into the event stack,
the event may be canceled at any time prior to the
time of occurrence. This is accomplished with a FORTRAN
call of the following form,

> CALL CANCEL(mHname,dummy,status)

where,

m is a character count of the number of
characters contained in the name of the
event $(1 \le n \le 5)$.

name is the name of the event routine (left
justified).

dummy is a dummy argument which is not used, but
must be included for compatibility with
calls to EVENT.

status is an integer variable which is used to
signal the status of the cancellation
operation to the user.

Upon return from the cancellation routine status con-
tains,

0 if the routine was found in the event
stack and was successfully canceled.

1   if the routine was not found in the stack

and no action was taken.

2   if the event stack was empty.

If an event is scheduled to occur more than once and a call to CANCEL is made, the entry which was first to occur is removed from the event stack.

System-defined events.

A number of events are defined and scheduled by the system. The user should be made aware of these events for the following reason. In some simulations various system actives are sometimes scheduled to occur at the same time as user-defined events. Most notable of these is the system routine which produces printed output. If a user's routine was scheduled to occur at the same time as the system's printing routine but at a lower priority, then the printed output would not reflect the state of the system after events scheduled at that time have occurred. Table 2.1.3-1 contains a list of the system routines, their scheduling priority, and the system-defined variable which controls their time of occurrence. Some of the routines listed are user-defined special purpose subroutines described in section 2.1.4.

Table 2.1.3-1. System-defined events.

Routine Name	Priority	Controlling Variable	Action
START[a]	100	TSTRT	User-supplied.
XPRNT	200	DTPR	Prints tabular output.
XPLOT	200	DTPL	Saves values of variables for plotting.
XCSIM	300	DT	CYCL1 is called if included by the user, the flows are computed, the state variables are updated, and CYCL2 is called if included by the user.
FINIS[a]	500	TEND	User-supplied.
HALT	512	TEND	Halts execution.

[a] User-supplied routines, scheduled by the system.

The system-defined routine HALT can be scheduled
by the user any time he desires the simulation
terminated. A choice of priorities for scheduling
events should be made with the above table in mind.
In most situations the user will desire to schedule
his events at a higher priority than the system events
(i.e., less than 100). As indicated in Table 2.1.3-1, a
number of system-defined events are scheduled accord-
ing to the values given the reserved system control
variables. If values for these variables are needed
by the system, but have not been set by the user,
default values will be supplied. A complete discussion
of the system-control variables and their use in
controlling simulations is contained in section 2.2.1.

Construction of event simulations.

The construction and operation of event-oriented
simulations is illustrated by considering a simple-
event simulation of a hypothetical population. The
simulation is not intended to be biologically realistic.
The processes to be considered are (i) births, (ii)
recruitment from the juvenile age class to the adult
age class, and (iii) deaths. The following variables
are the variables of interest in the simulation.

NJM - No. of juvenile males.

NJF - No. of juvenile females.

NJ - Total no. of juveniles.

NAM - No. of adult males.

NAF - No. of adult females.

NA - Total no. of adults.

N   - Total population.

Each of the processes or events in the simulation consist of two sets of computations. These computations are (i) computations which reflect changes in the variables of interest due to the processes being simulated and (ii) computations which determine the time at which an event will occur. The particular equations used to compute these two quantities embody the assumptions about the processes involved in the population. For the sake of clarity in this example we will assume a very simple description for the processes influencing the population dynamics.

(1) Births are assumed only to occur from the 90th to the 120th day of each year. We assume that 80% of all female adults have offspring during this time interval and that the number of offspring per female occurs in the following proportions.

No. of offspring/birth	Percent occurrence
1	5%
2	80%
3	10%
4	5%

We further assume that males are born as often as females. Therefore during the 30 days of natality the number of birth events which occur on the average is 0.8 * NAF. Hence the average time between births is 30./(0.8 * NAF). The standard deviation of the time between births is assumed to be 10% of the mean.

(2) Recruitment from the juvenile age class to the adult age class is based on the assumption that the mean time required for a juvenile to mature is 365 days with a standard deviation of 20 days.

(3) Deaths are assumed to occur according to the following graph of mean adult lifespan vs. total population size.

The standard deviation of the mean adult lifespan is assumed to be 10% of the mean. Table 2.1.3-2 lists the events required for this simulation and the actions performed by each event.

Table 2.1.3-2. Population simulation events.

Event	Name	Computations
(1) Birth	BIRTH	(a) No. of offspring/birth
		(b) Sex of each offspring
		(c) Adjust population size
		(d) Schedule time of maturation
		(e) Schedule time of next birth
(2) Male/female maturation	RCRTM/ RCRTF	(a) Adjust age class sizes.
		(b) Schedule time of death
(3) Male/female death	DETHM/ DETHF	(a) Adjust population size

A complete listing of the simulation and the results are contained in example 2.1.3-1. Since the simulation contains stochastic elements, the results shown in the output represent only one of the many relaizations of the simulation which would be required for an exhaustive analysis of the model. Whenever a variable in a simulation is defined stochastically, the value of the variable at any point in the simulation is obtained by the random sampling of a value from the indicated distribution function. Therefore any one run of the simulation represents only one possible realization of the system which is being modeled. If the statistical properties of the variables of interest are desired, a number of runs using different random number sequences would be required.

Example 2.1.3-1. A sample simulation illustrating events.

```
                  STORAGE. NJM.NJF.NJ.NAM.NAF.NA.N.YEAR
                  UNIFORM. FRCT
                  NORMAL. TSMP
                      EVENT BIRTH
                      YEAR=TIME/365.
                  C...BIRTH EVENT.  DETERMINE THE NUMBER OF OFFSPRING.
                      F=FRCT(0.,1.)
                      NR=1
                      IF(F.LE.0.05) GO TO 5
                      NB=2
                      IF(F.LE.0.85) GO TO 5
                      NR=3
                      IF(F.LE.0.95) GO TO 5
                      NR=4
                  C...INCREMENT THE POPULATION VARIABLES AND SCHEDULE RECRUITMENT.
                    5 DO 20 I=1,NB
                  C...SAMPLE THE TIME OF RECRUITMENT OF THE OFFSPRING.
                   10 TRC=TSMP(365.,20.)
                      IF(TRC.LE.0.) GO TO 10
                  C...INCREMENT THE TOTAL POPULATION SIZE AND NO. OF JUVENILES.
                      NJ=NJ+1
                      N=N+1
                  C...DETERMINE THE SEX OF THE OFFSPRING.
                      R=FRCT(0.,1.)
                      IF(R.GT.0.5) GO TO 15
                      NJM=NJM+1
                      CALL EVENT(5HRCRTM,TIME+TRC,20)
                      GO TO 20
                   15 NJF=NJF+1
                      CALL EVENT(5HRCRTF,TIME+TRC,20)
                   20 CONTINUE
                  C...SCHEDULE THE TIME TO THE NEXT BIRTH.
                      TMB=30./(0.8*NAF)
                      TSB=0.1*TMB
                   25 TB=TSMP(TMB,TSB)
                      IF (TB.LE.0.) GO TO 25
                      TY=AMOD(TIME,365.)
                      IF(TY+TB.GT.120.) TB=TB+335.
                      CALL EVENT(5HBIRTH,TIME+TB,20)
                      RETURN
                      END
                      EVENT RCRTM
                      YEAR=TIME/365.
                  C...EVENT OF THE RECRUITMENT OF A MALE JUVENILE.
                      NJM=NJM-1
                      NAM=NAM+1
                      NJ=NJ-1
                      NA=NA+1
                  C...SAMPLE THE LIFESPAN OF THE ADULT AND SCHEDULE THE DEATH.
                      TML=ALINT2(N,IFLG,100,1095,500,365,1000,36)
                      TSL=0.1*TML
                    5 TD=TSMP(TML,TSL)
                      IF (TD.LE.0.) GO TO 5
                      CALL EVENT(5HDETHM,TIME+TD,20)
                      RETURN
                      END
                      EVENT RCRTF
                      YEAR=TIME/365.
                  C...EVENT OF THE RECRUITMENT OF A FEMALE JUVENILE.
                      NJF=NJF-1
                      NAF=NAF+1
                      NJ=NJ-1
                      NA=NA+1
                  C...SAMPLE THE LIFESPAN OF THE ADULT AND SCHEDULE THE DEATH.
                      TML=ALINT2(N,IFLG,100,1095,500,365,1000,36)
                    5 CALL EVENT(4HHALT,TIME,1)
                      RETURN
                      END
                      EVENT DETHF
                      YEAR=TIME/365.
                  C...DEATH EVENT.  IF POPULATION GOES TO ZERO THE SIMULATION IS HALTED.
                      IF(N.LT.1) GO TO 5
                      N=N-1
                      NA=NA-1
                      NAF=NAF-1
                      RETURN
                    5 CALL EVENT(4HHALT,TIME,1)
                      RETURN
                      END
                      SUBROUTINE START
```
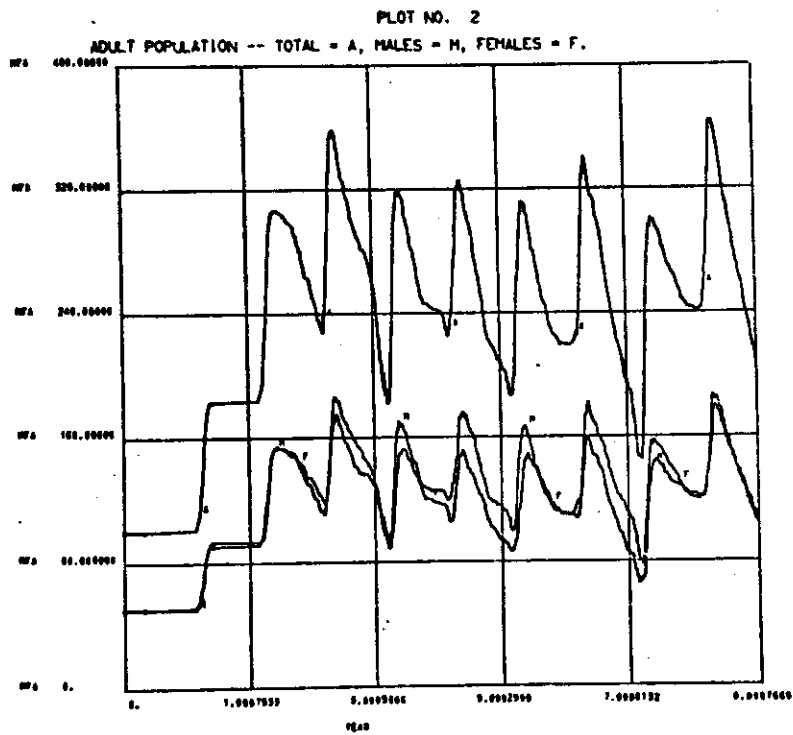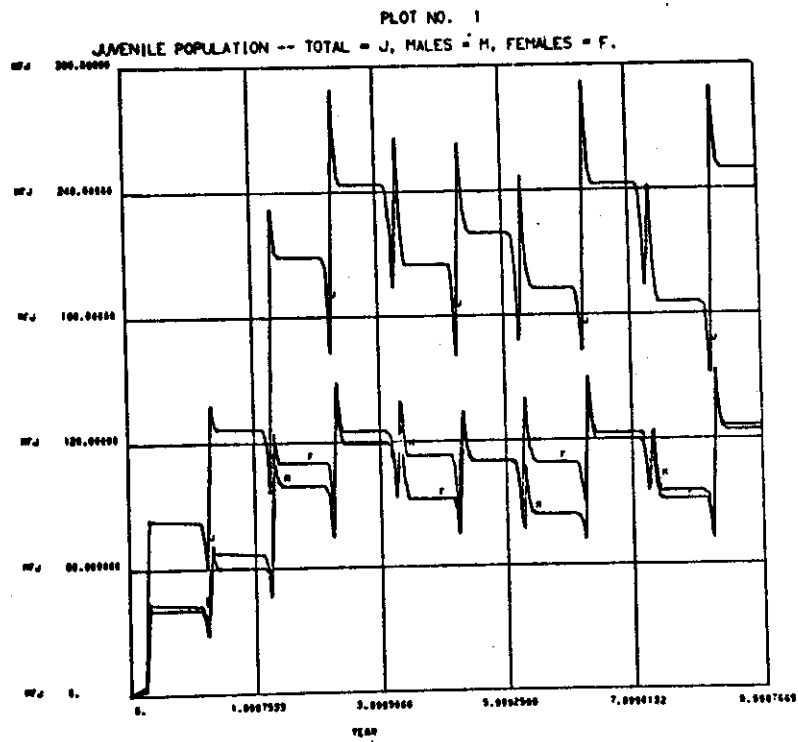
```
            TSL=0.1*TML
        5 TD=TSMP(TML,TSL)
          IF (TD.LE.0.) GO TO 5
          CALL EVENT(5HDETHF,TIME+TD,20)
          RETURN
          END
          EVENT DETHM
          YEAR=TIME/365.
C...DEATH EVENT.  IF POPULATION GOES TO ZERO THE SIMULATION IS HALTED.
          IF(N.LT.1) GO TO 5
          N=N-1
          NA=NA-1
          NAM=NAM-1
          RETURN
C...ASSUMING AN INITIAL POPULATION OF YOUNG ADULTS ONLY, SCHEDULE THEIR
C       DEATHS.
          DO 5 I=1,NAM
          TML=ALINT2(N,IFLG,100,1095,500,365,1000,36)
          TSL=0.1*TML
        4 TD=TSMP(TML,TSL)
          IF (TD.LE.0.) GO TO 4
        5 CALL EVENT(5HDETHM,TIME+TD,20)
          DO 10 I=1,NAF
          TML=ALINT2(N,IFLG,100,1095,500,365,1000,36)
          TSL=0.1*TML
        9 TD=TSMP(TML,TSL)
          IF (TD.LE.0.) GO TO 9
       10 CALL EVENT(5HDETHF,TIME+TD,20)
C...SCHEDULE THE FIRST BIRTH.
          CALL EVENT(5HBIRTH,90.,20)
          RETURN
          END
```
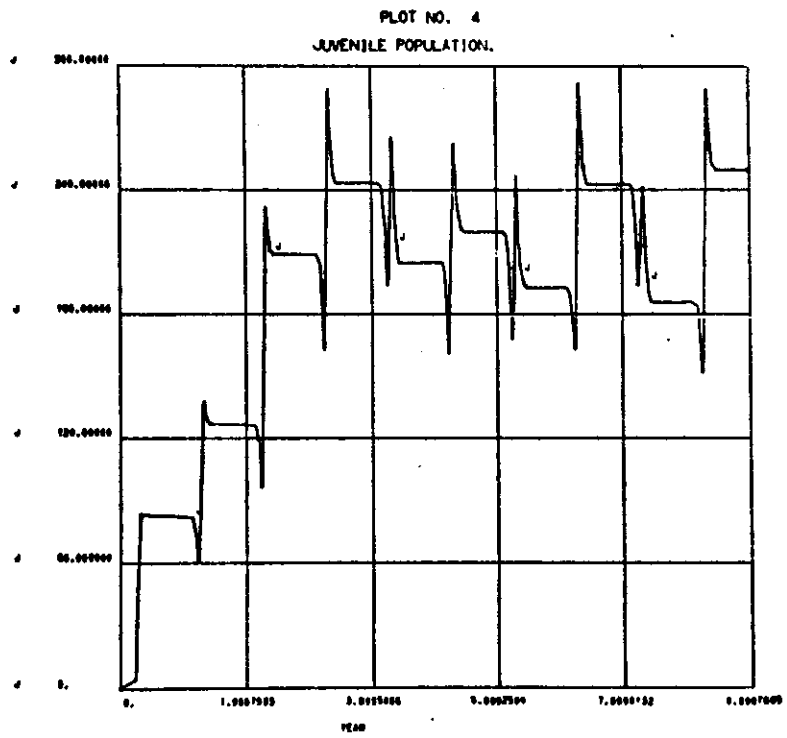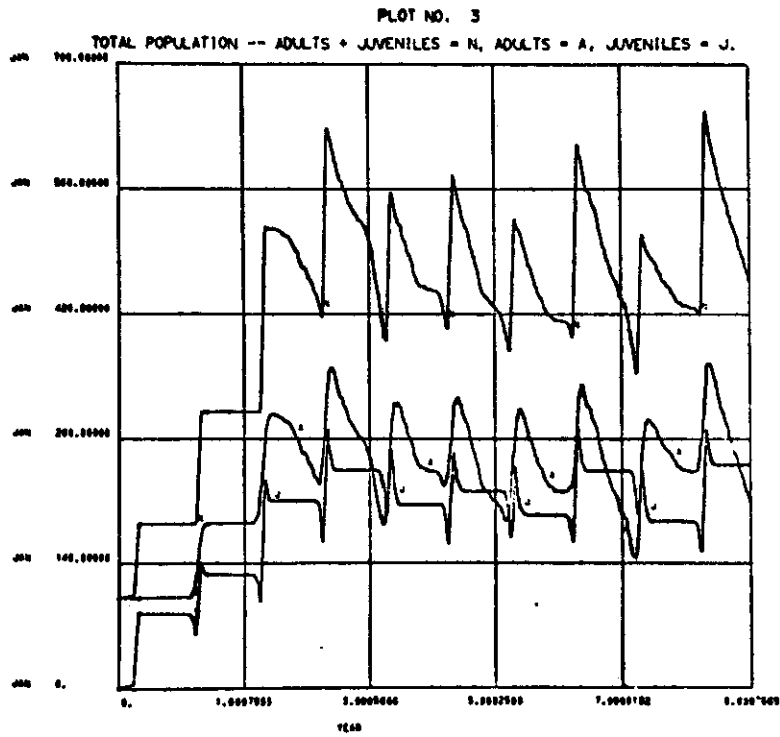
7_89 end-of-record separator

```
NJM=0 $ NJF=0 $ NJ=0 $ NAM=50 $ NAF=50 $ NA=100 $ N=100 $
YEAR=0. $ TSTRT=0. $ TEND=3650. $ DTPL=4.5625 $
TITLE. JUVENILE POPULATION -- TOTAL = J, MALES = M, FEMALES = F.
PLOT. (NJM=M,NJF=F,NJ=J)/YEAR
TITLE. ADULT POPULATION -- TOTAL = A, MALES = M, FEMALES = F.
PLOT. (NAM=M,NAF=F,NA=A)/YEAR
TITLE. TOTAL POPULATION -- ADULTS + JUVENILES = N, ADULTS = A, JUVENILES = J.
PLOT. (NJ=J,NA=A,N=N)/YEAR
TITLE.                        JUVENILE POPULATION.
PLOT. (NJ=J)/YEAR
TITLE.                        ADULT POPULATION.
PLOT. (NA=A)/YEAR
TITLE.                        TOTAL POPULATION.
PLOT. (N=N)/YEAR
FILM.
```

GRAPHICAL SIMULATION RESULTS   07/23/73   10.11.52.

GRAPH NO.	GROUP	GROUP RANGE DECLARATION	DEPENDENT VARIABLE(S)	PLOTTED CHARACTER	INDEPENDENT VARIABLE	INDEPENDENT VARIABLE RANGE DECLARATION
1	1		NJM NJF NJ	M F J	YEAR	
2	1		NAM NAF NA	M F A	YEAR	
3	1		NJ NA N	J A N	YEAR	
4	1		NJ	J	YEAR	
5	1		NA	A	YEAR	
6	1		N	N	YEAR	

PLOT NO. 1

JUVENILE POPULATION -- TOTAL = J, MALES = M, FEMALES = F.

PLOT NO. 2

ADULT POPULATION -- TOTAL = A, MALES = M, FEMALES = F.

PLOT NO. 3

TOTAL POPULATION -- ADULTS + JUVENILES = N, ADULTS = A, JUVENILES = J.

PLOT NO. 4

JUVENILE POPULATION.

PLOT NO. 5
ADULT POPULATION.

PLOT NO. 6
TOTAL POPULATION.

2.1.4 Subprograms.

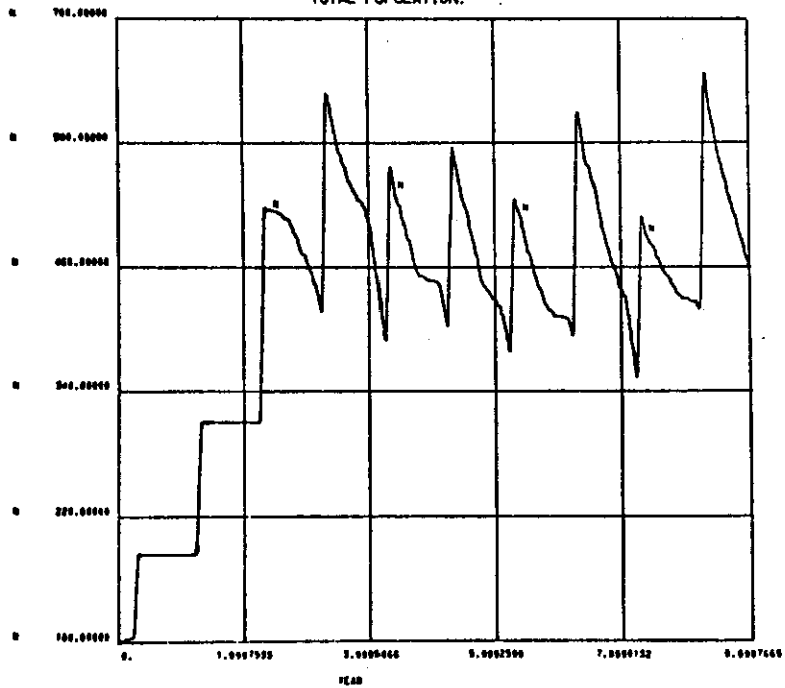FORTRAN subroutines and functions may be supplied by the programmer. Subroutines and functions may appear anywhere in the source section provided they do not appeear within the intended range of a flow definition label. Parameter declaration statements may not be included within the text of a subprogram. All reserved system control variables (cf. Table 2.1-1) and all user-defined variables and stochastic functions declared in parameter declaration statements should be considered globally defined and are accessible within every user-supplied subprogram. *These variables are in a common block inserted by SIMCOMP into each of these routines.* The format of user-supplied subprograms conforms to FORTRAN specifications for subroutines and functions. A user-supplied subprogram can be called from within any flow definition or other subprogram or event.

Certain special purpose subroutines can be supplied by the user which will be called by the SIMCOMP executive timing routine at predetermined times in the simulation. These reserved routine names are listed in Table 2.1.4-1. Computations which are to be performed at the specified times are included in a FORTRAN subroutine appropriately named. Since special-purpose subroutines are called by the executive routine, argument lists are not allowed. A flow chart of the execution sequence in a simulation containing flow definitions is given in Fig. 2.1.4-1.

Table 2.1.4-1.   Reserved subroutine names

Subroutine Name	Use
START	Called after parameter values have been set by data assignment statements in the data section just prior to the start of execution, TIME = TSTRT.
CYCL1[a]	Called just prior to the computation of flows at each time step.
CYCL2[a]	Called after the flows have been computed and the state variables have been updated, but prior to any printing or storing of values for plotting at each time step.
FINIS	Called at the end of simulation, TIME = TEND.

[a] Called only if flow definitions are present.

Fig. 2.1.4-1. Flow execution sequence.

Example 2.1.4-1. An example of a user-defined subroutine. The graph was
reproduced from a printer plot generated by SIMCOMP.

```
STORAGE. V(5),P1,P2
(I=101,105-263).
       CALL VCALC
       J=I-100
       FLOW=V(J)*P1/DT
       SUBROUTINE VCALC
       DO 10 I=1,5
       J=I+100
   10  V(I)=X(J)*P2
       RETURN
       END
```

7_89 *end-of-record separator*

```
X(101)=10.,12.,15.,8.,20. $ X(263)=1000. $ P1=0.0001 $ P2=2.363 $
TSTRT=0. $ TEND=100. $ DT=2. $
PLOT. (X(101)=1),(X(102)=2),(X(103)=3),(X(104)=4),(X(105)=5)
```

Note that the variables in the STORAGE. statement
are globally defined and are available for use in the
subroutine and in the flows.

Example 2.1.4-2. An example of a user-supplied function.

```
STORAGE. A.T1.T2.PHI.C1.C2
        FUNCTION COSX(T)
        COSX=A*COS((T-T2)/3.14+PHI)
        RETURN
        END
        SUBROUTINE CYCL1
        C1=X(9)*COSX(TIME-T1)
        C2=0.1*C1
        RETURN
        END
(1-9). FLOW=C1
(2-9). FLOW=C2/X(2)
```

789   end-of-record separator

```
A=0.01 $ T1=20. $ T2=10. $ PHI=3.14 $ C1=0. $ C2=0. $
TSTRT=0. $ TEND=6.28 $ DT=0.0628 $
X(1)=100. $ X(2)=100. $ X(9)=1. $
PLOT. (X(9)=9)
```

PLOT NO. 1

2.1.5 Utility routines.

A library of subroutines and functions is available which is accessed and made available to the user by SIMCOMP. Whenever the user includes a call to one or more of the utility routines and does not supply a subprogram by the same name, the utility library is accessed and the called routine or routines are loaded. The routines currently available include the functions reported by Parton and Innis (1972). The calling sequences for their functions are compatible with the FORTRAN listings provided therein. The utility routines listed in Table 2.1.5-1 are also available, and a description of their use follows.

Table 2.1.5-1. Utility routines.

Name	Purpose
ALINT1	Linear interpolation of data whose independent variable is regularly spaced.
ALINT2	Same as ALINT1 for unevenly spaced data points.
FLOWV	Returns the most recently computed value for a particular flow.
XSTATS	Statistical sampling package
PUNCHD	Produces a punched deck of data in SIMCOMP acceptable format.

Linear interpolation.

Many times the most desirable way to specify a function is by a table which is linearly interpolated. The FORTRAN callable functions ALINT1 and ALINT2 will interpolate a table of values producing a value of the dependent variable for any given value of the independent variable.

Mathematically, the operation of ALINT1 and ALINT2 is described.

Given:

x[3] - the value of the independent variable at which point a linearly interpolated value is to be computed.

n - the number of pairs of values in the interpolation table.

x_j - the jth value of the independent variable in the table to be interpolated.

y_j - the jth value of the dependent variable in the table corresponding to x_j.

In order for the linear interpolation routine ALINT2 to operate efficiently, the values of the independent variable must be in ascending order,

$$x_j \leq x_{j+1} \quad \text{for} \quad j = 1, 2, \ldots, n-1.$$

The functions ALINT1 and ALINT2 compute the linearly interpolated value as follows:

---

[3] The use of lower case x in this section should not be confused with the state variables X.

$$\text{if } x \leq x_1 \text{ then ALINT} = y_1$$

$$\text{if } x \geq x_n \text{ then ALINT} = y_n$$

$$\text{if } x_j \leq x < x_{j+1} \text{ for } j = 1, 2, \ldots, n - 1$$

then

$$\text{ALINT} = y_j + \frac{y_{j+1} - y_j}{x_{j+1} - x_j} \cdot (x - x_j).$$

Equal interval data.

The utility function ALINT1 can be used whenever the values of the independent variable in the interpolation table are equally spaced. An example of such a table and its graph is presented in Fig. 2.1.5-1. The linearly interpolated value for x = 0.25 is y(x) = 0.375. Note in the graph that for values of the independent variable outside the range of definition of the table, that is less than 0.0 and greater than 0.5, the value of the dependent variable is assumed to be equal to the value of the function at its tabular end point. Therefore x = 0.55 produces y(x) = 0.05. Whenever extrapolation occurs, this is the action taken by ALINT1 and ALINT2.

x_j	Y_j
0.0	.1
0.1	.3
0.2	.4
0.3	.35
0.4	.25
0.5	.05

Fig. 2.1.5-1. Sample tabular function and graph containing equal interval data.

The second calling sequence to ALINT1 is,

$$V = ALINT1(R,IFLG,RS,RD,Y1,Y2, \ldots , YN)$$

where R, IFLG, RS, and RD are as defined in the first type of call. The constants or variables Y1 through YN correspond to the values in the array RT. The quantities Y1 through YN correspond to the values y_1, y_2, $\ldots$, y_n in the mathematical description of the interpolation algorithm. These quantities may be either integer or real valued, but again it is recommended that real values be used. There must be at least two or more entries of the dependent variable in the call(i.e., n must be greater than or equal to two). Using the second method of calling, the previous example would be programmed as follows.

Example 2.1.5-2. Use of ALINT1, type 2 call.

```
(1-2). FLOW=ALINT1(X(2),IDUM,0.,0.1,0.1,0.3,0.4,0.35,0.25,0.05)
78,    end-of-record separator
X=100.,0. $ TSTRT=0. $ TEND=10. $ DT=0.1 $
PLOT. (X(2)=2)
```

Unequally spaced data.

The utility function ALINT2 must be used whenever the values of the independent variable in the interpolation table are unequally spaced. In this case the values of the independent variable at each of the tabular points must be supplied in the function call. An example of an interpolation table which contains unequally spaced data is presented in Fig. 2.1.5-2.

The first type of FORTRAN calling sequence to
ALINT1 is,

$$V = ALINT1(R,IFLG,RS,RD,RT)$$

where, R   is the value of the independent variable
at which point a linearly interpolated value
is desired.

RS   is the starting value of the independent
variable in the table.

RD   is the increment between successive values
of the independent variable in the table.

RT   is an array containing the tabular values
of the dependent variable.

The above four quantities can be either integer- or real-
valued variables or constants. It is recommended that
only real values are used so that internal conversion
is not required. The array RT must be declared in
STORAGE., and the declared size of the array must be
one word longer than the number of values in the array.
The last location in the array must not be set to any
value either in the data section or within the source
section. By use of this convention, ALINT1 is able to
determine n, the number of values in the table. Upon
return from the function,

ALINT1   is the interpolated value.

IFLG   is an integer variable set by the function
as a flag to the user,

IFLG = 0   for normal interpolation.

IFLG = 1   if extrapolation occurred.

For the sample data in Fig. 2.1.5-1 the above calling

sequence is illustrated.

Example 2.1.5-1. Use of ALINT1, type 1 call.

```
STORAGE. Y(7)
(1-2). FLOW=ALINT1(X(2),IDUM,0.,0.1,Y)
⁷8₉    end-of-record separator
Y=0.1,0.3,0.4,0.35,0.25,0.05 $ X=100.,0. $ TSTRT=0. $ TEND=10. $ DT=0.1 $
PLOT. (X(2)=2)
```

X_i	Y_i
0	1
10	.7
20	.5
35	.35
50	.25
70	.2

Fig. 2.1.5-2. Sample tabular function and graph containing unequally spaced data.

The first type of FORTRAN calling sequence to ALINT2 is,

$$V = ALINT2(R,IFLG,TABLE)$$

where R and IFLG are as described in the calls to ALINT1. The variable TABLE must be a two-dimensional array declared in STORAGE. and is dimensioned TABLE (2,n+1) where n is the number of pairs of values in the interpolation table. The entries TABLE (1,n+1) and TABLE (2,n+1) must not be given any values in the data section or in the source program. The array TABLE can be either an integer-valued or real-valued array, but it is recommended that a real-valued array be used. The location of the values of the interpolation table in the array TABLE follows,

$$TABLE\ (1,1) = x_1, \quad TABLE\ (2,1) = y_1$$

$$TABLE\ (1,2) = x_2, \quad TABLE\ (2,2) = y_2$$

$$\vdots \qquad\qquad\qquad \vdots$$

$$TABLE\ (1,n) = x_n, \quad TABLE\ (2,n) = y_n,$$

with TABLE (1,n+1) and TABLE (2,n+1) not given any

values. The following example is derived from the

data given in Fig. 2.1.5-2.

Example 2.1.5-3. Use of ALINT2, type 1 call.

```
STORAGE. TT(2,7)
(8-80). TV=0.8*X(80)
       TF=ALINT2(TV,IFL,TT)
       FLOW=TV*TF
       IF(TV.LT.0.) FLOW=0.
78₉      end-of-record separator
TT=0.,1.,10.,0.7,20.,0.5,35.,0.35,50.,0.25,70.,0.2 $
TSTRT=0. $ TEND=10. $ DT=0.1 $ X(8)=500. $ X(80)=50. $
PLOT. (X(8)),(X(80))
```


PLOT NO. 1

The second type of calling sequence to ALINT2 is,

V = ALINT2(R,IFLG,R1,Y1,R2,Y2, ... ,RN,YN)

where R and IFLG are as described in the calls to ALINT2
ALINT1. The constants or variables R1, Y1 through RN,
YN correspond to the entries in TABLE in the first type
of call to ALINT2. Referring to the mathematical
description of linear interpolation the arguments R1,
Y1 through RN, YN correspond to,

$$R1 = x_1, \quad Y1 = y_1$$
$$R2 = x_2, \quad Y2 = y_2$$
$$\vdots \qquad \vdots$$
$$RN = x_n, \quad YN = y_n.$$

Note that in both types of calls to ALINT2 the values
of the independent variable in the interpolation table
must be specified in ascending order. Using the second
method of calling ALINT2, the sample data contained in
Fig. 2.1.5-2 would be linearly interpolated as in the
following example.

Example 2.1.5-4. Use of ALINT2, type 2 call.

```
(8-80). TV=0.8*X(80)
        TF=ALINT2(TV,IFL,0,1,10,0.7,20,0.5,35,0.35,50,0.25,70,0.2)
        FLOW=TV*TF
        IF(TV.LT.0.) FLOW=0.
⁷⁸₉    end-of-record separator
TSTRT=0. $ TEND=10. $ DT=0.1 $ X(8)=500. $ X(80)=50. $
PLOT. (X(8)),(X(80))
```

Step functions.

The special utility function ALINT2 can be used to
generate step functions. A step function $f(x)$ is in

general defined by,

$$f(x) = \begin{cases} a & \text{if } x \le x_s \\ b & \text{if } x > x_s. \end{cases}$$

By allowing two consecutive entries of the independent variable in the interpolation table to assume the same value, ALINT2 will generate a step. Fig. 2.1.5-3 contains the equation of a step function, its graph, and the interpolation table used for its generation. Using the second form of call to ALINT2, the step function illustrated in Fig. 2.1.5-3 would be programmed as in the following definition of a flow.

Fig. 2.1.5-3. Sample step function.

Example 2.1.5-5. Use of ALINT2 as a step function.

```
STORAGE. RAIN
(3-1). RAIN=5.106*EXP(-0.00125*(TIME-50.)**2)
      FLOW=RAIN
(1-2). FLOW=0.1*X(1)*ALINT2(RAIN,IF,3.,0.,3.,1.)
789      end-of-record separator
X=0.,0.,100. $ TSTRT=0. $ TEND=100. $ RAIN=0.22434 $ DT=1. $
PLOT. (X(1)),(X(2)),(RAIN)
```

PLOT NO. 1

```
A    80.0000
B    200.000
C    6.00001

A    64.0000
B    160.000
C    4.80000

A    48.0000
B    120.000
C    3.60000

A    32.0000
B    80.0000
C    2.40000

A    16.0000
B    40.0000
C    1.20000

A    0.
B    0.
C    0.

     0.        20.0000      40.0000      60.0000      80.0000      100.000

                              TIME
```

Retrieving values of flows.

The value which is computed for a flow is not directly available to any part of a source program except within the range of the particular flow. Sometimes it is desirable to acquire the value of a particular flow or flows for use in the computation at some later time. The usual situation is when the value of one flow depends upon the value of some previously computed flow. While this situation is physically impossible (Innis 1972), it may be a very useful procedure. The special utility subroutine

FLOWV is used to access the value of a previously
computed flow given the source and destination
compartment indices of the desired flow. The FORTRAN
calling sequence to the routine FLOWV is,

            CALL FLOWV(I,J,VALUE,IFLAG)

where   I   is the source compartment index.

             J   is the destination compartment index.

             VALUE is the most recently computed value

                     of the flow.

             IFLAG is a flag which signals various

                     conditions to the user.

Table 2.1.5-2 summarizes the values which IFLAG can
attain and their meanings. Any attempt to subsequently
use the variable VALUE in the computation while IFLAG
returning any of the values 1 through 5 will result
in an arithmetic-mode error and the simulation will be
abnormally terminated. By checking the value of IFLAG
prior to using VALUE in the computation, an abnormal
termination can be avoided. Refer to section 2.3 to
determine what arithmetic operations can produce any
of the conditions 2 through 4. The quantities I and
J may be integer constants or variables. VALUE must
be a real-valued variable. The following flow chart
and example illustrate the use of FLOWV.

Table 2.1.5-2.   Values for IFLAG on return from FLOWV.

IFLAG	VALUE	Description
0	--	Current value returned OK.
1	Indefinite	Flow exists, but was not assigned a value by the user.
2	+Infinite	Flow exists, value is positive infinite.
3	-Infinite	Flow exists, value is negative infinite.
4	Indefinite	Flow exists, value is indefinite.
5	Indefinite	Flow was not defined in the simulation.

Example 2.1.5-6.   Simulation containing a call to FLOWV.

```
STORAGE. R(3)
(5-4). FLOW=10.*SIN(TIME*3.14159/365.)
(4-I=1,3). CALL FLOWV(5,4,FL,ICHK)
      FLOW=0.
      IF(ICHK.NE.0) GO TO 10
      FLOW=R(I)*FL
   10 CONTINUE
```

789   *end-of-record separator*

```
X=4*0.,5000.  $ R=0.1,0.3,0.5  $
TSTRT=0.  $ TEND=365.  $ DT=1.  $
PLOT.  (X(5)=5),(X(4)=4)
PLOT.  (X(1)=1,X(2)=2,X(3)=3)
```

PLOT NO. 1

PLOT NO. 2

Statistical-sampling package.

        Often the outputs of a simulation experiment are statistical measurements. Statistical measures of simulated variables through time are often the principle variables of interest in event simulations. Such quantities as the yearly average population size, the mean time to maturation, and the probability of reaching a given age are typical measures of performance of an event-oriented simulation of population dynamics. Normally, statements must be scattered throughout the program to gather such statistics. Writing the statements necessary for gathering such quantities as sums and sums of squares is a task to be avoided because it clutters the logic of simulation with statements whose only function is the collection of output information.

        The statistical sampling package XSTATS provides a number of subroutines which simplify the gathering and reporting of such information. The statistical sampling routines allow the use of either of two sampling strategies. These are (i) discrete sampling of variables and (ii) time-weighted sampling of variables. The mathematical description of each of these methods is presented in Table 2.1.5-3 given the following definitions.

        n   -   the number of samples.

x_i [4/] — the value of the ith sample.

t_i — the time at which the value changed from x_i to x_{i+1}.

t_0 — the time at which sampling started.

$\Delta t_i = t_i - t_{i-1}$ — the length of time which the ith sample had the value x_i.

Table 2.1.5-3. Statistical sampling computational methods.

Statistic	Discrete Sampling	Time-weighted Sampling
Number or total time	n	$\Sigma \Delta t_i$
Sum	Σx_i	$\Sigma x_i \cdot \Delta t_i$
Sum of squares	Σx_i^2	$\Sigma x_i^2 \cdot \Delta t_i$
Mean	$\dfrac{\Sigma x_i}{n}$	$\dfrac{\Sigma x_i \cdot \Delta t_i}{\Sigma \Delta t_i}$
Mean square	$\dfrac{\Sigma x_i^2}{n}$	$\dfrac{\Sigma x_i^2 \cdot \Delta t_i}{\Sigma \Delta t_i}$
Variance	$\dfrac{\Sigma x_i^2}{n} - \left(\dfrac{\Sigma x_i}{n}\right)^2$	$\dfrac{\Sigma x_i^2 \cdot \Delta t_i}{\Sigma \Delta t_i} - \left(\dfrac{\Sigma x_i \cdot \Delta t_i}{\Sigma \Delta t_i}\right)^2$
Standard deviation	$\sqrt{\text{variance}}$	$\sqrt{\text{variance}}$
Maximum	largest x_i	largest x_i
Minimum	smallest x_i	smallest x_i

Discrete sampling.

As an example of discrete sampling consider the

following table of values which might have been the

_____

[4/] The use of lower case x in this section should not be confused with the state variables X.

times to maturation of individuals in a simulated population.

individual	1	2	3	4	5	6
time to maturity	1.8	1.7	2.1	0.9	1.6	1.7

Table 2.1.5-4 displays the values of each of the statistics using discrete sampling.

Table 2.1.5-4. Example of discrete sampling.

Statistic	Value
Number	6.0
Sum	9.8
Sum of squares	16.8
Mean	1.63
Mean square	2.8
Variance	0.132
Standard deviation	0.364
Maximum	2.1
Minimum	0.9

Time-weighted sampling.

As an example of time-weighted sampling consider the following table of values and their graph through time. The graph might be the graph of the size of a

population through time. Remember that the value t_i
is the time at which the sampled variable changed from
x_i to x_{i+1}.

t_i	1.4	4	7.4	11.9	16.5	19.6	22	22.8	24
Δt_i	1.4	2.6	3.4	4.5	4.6	3.1	2.4	0.8	1.2
x_i	0	1	3	2	5	6	4	3	2

Using time-weighted statistics, we are not just interested
in quantities such as the average of each of the sizes
the population takes on, but in the more meaningful
quantities such as the time-weighted average. That is,
the average over the 24-unit time interval is computed
by averaging the sizes the population assumes weighted
by the time spent in each population size. Table 2.1.5-5
displays the values of each of the statistics using
time-weighted sampling for the above example.

Table 2.1.5-5. Example of time-weighted sampling.

Statistic	Value
Total time	24.0
Sum	77.8
Sum of squares	328.2
Mean	3.242
Mean square	13.675
Variance	3.167
Standard deviation	1.779
Maximum	6.0
Minimum	0.0

Statistical package calling sequences.

The statistical sampling package XSTATS contains two entry points for sampling values of variables. The FORTRAN-calling sequences for sampling the value of a variable are

STORAGE.   VAR
    .
    .

    CALL SAMPLE(VAR)

or   CALL SAMPLE(VAR,TIME).

The first argument in the call must be a real-valued variable and should be declared in STORAGE. The first form of the call is used for discrete sampling. Each time a new value for the sampled variable is computed,

a call to sample should be made. The second form of
call is used for time-weighted sampling. A call to
sample should be made for each value the variable
assumes when TIME is equal to the final end point of
the interval over which the value holds. Calls to
SAMPLE for a particular variable can not be mixed
between the two forms of call. If an attempt is made
to sample a variable in both discrete and time-weighted
modes, a diagnostic is issued and subsequent calls
referencing this variable are ignored. A maximum of
30 different variables can be sampled within a simula-
tion. Upon attempting to sample more than 30 variables,
a diagnostic is issued and attempts to sample in excess
of the first 30 variables are ignored.

The sampling sequence can be reinitialized to
begin anew at a point in the simulation by calls of
the following form,

$$\text{CALL} \quad \text{RESET(VAR)}$$

or $$\quad \text{CALL} \quad \text{RESET(VAR,TIME)}.$$

When RESET is called, each of the statistics for the
named variable are reset to the appropriate initial
values. In the discrete sampling case quantities such
as n, Σx and Σx^2 are initialized to zero. For time-
weighted sampling the initial time of sampling t_0 is
set to the current value of time in addition to the
other required initializations. If RESET is not called

prior to any time-weighted calls to SAMPLE, the initial value assumed for t_0 is TSTRT, or the time of the first event if TSTRT is not given a value in an event-only simulation.

The statistics gathered by calls to sample will be printed by executing a call of the following form,

CALL REPORT(VAR)

A report of all statistics will be printed in the output as illustrated in the example at the end of this section automatically. If the variable VAR is declared in STORAGE., the output will be labeled with the name of the variable. Otherwise the output is labeled by an integer enclosed in asterisks (*) which represents the actual core location of the variable.

The values of any of the statistics for any sampled variable can be assessed during the simulation by the FORTRAN function calls presented in Table 2.1.5-6. If the requested variable has not been sampled prior to the function call, an indefinite result is returned.

Table 2.1.5-6.   Statistical function calls.

Sample Calling Sequence	Value Returned
V = COUNT(VAR)	n or $\Sigma\Delta t_i$
V = SUM(VAR)	sum
V = SUMSQ(VAR)	sum of squares
V = AVERAGE(VAR)	mean
V = RMEANSQ(VAR)	mean - square
V = VARIANC(VAR)	variance
V = STDEV(VAR)	standard deviation
V = RMAX(VAR)	largest value
V = RMIN(VAR)	smallest value

Example 2.1.5-7.   Illustration of calls to the statistical package.

```
STORAGE. POP,BDEL,DDEL,PMEAN
NORMAL. PNORM
      EVENT BIRTH
C...EACH TIME THIS EVENT IS CALLED, A BIRTH OCCURES.  THE TIME BETWEEN
C     BIRTHS IS ASSUMED TO BE NORMALLY DISTRIBUTED WITH THE MEAN AND
C     STANDARD DEVIATION A FUNCTION OF TIME GIVEN BY AN INTERPOLATION
C     TABLE.
C
      IF(POP.LE.0.) RETURN
C...SAMPLE THE POPULATION SIZE.
      CALL SAMPLE(POP,TIME)
C...INCREMENT THE POPULATION SIZE.
      POP=POP+1.
C...SCHEDULE THE TIME OF NEXT BIRTH.
      PMEAN=ALINT2(TIME,ICHK,0.,4.,45.,2.5,70.,1.,90.,0.57,120.,0.5,
     -      180.,0.73,210.,1.25,270.,3.3,365.,4.)
      PSTDV=0.1*PMEAN
    5 BDEL=PNORM(PMEAN,PSTDV)
C...THE NORMAL DISTRIBUTION IS TRUNCATED AT ZERO.
      IF(BDEL.LE.0.) GO TO 5
      CALL EVENT(5HBIRTH,TIME+BDEL,1)
C...SAMPLE THE TIME BETWEEN BIRTHS.
      CALL SAMPLE(BDEL)
      RETURN
      END
      EVENT DEATH
C...EACH TIME THIS EVENT IS CALLED A DEATH OCCURES.  THE TIME BETWEEN
C     DEATHS IS ASSUMED TO BE NORMALLY DISTRIBUTED WITH A MEAN OF 1.0 AND
C     STANDARD DEVIATION OF 0.1.
C
```

```
        IF(POP.LE.0.) RETURN
C...SAMPLE THE POPULATION SIZE.
        CALL SAMPLE(POP,TIME)
C...DECREMENT THE POPULATION SIZE.
        POP=POP-1.
C...SCHEDULE THE TIME OF NEXT DEATH.
      5 DDEL=PNORM(1.,0.1)
C...THE NORMAL DISTRIBUTION IS TRUNCATED AT ZERO.
        IF(DDEL.LE.0.) GO TO 5
        CALL EVENT(5HDEATH,TIME+DDEL,1)
C...SAMPLE THE TIME BETWEEN DEATHS.
        CALL SAMPLE(DDEL)
        RETURN
        END
        EVENT STOP
C...REPORT STATISTICS.
        CALL REPORT(POP)
        CALL REPORT(BDEL)
        CALL REPORT(DDEL)
        RETURN
        END
    78g     end-of-record separator

TSTRT=0. $ TEND=365. $ POP=100. $ PMEAN=4. $
EVENT. BIRTH,0.,1
EVENT. DEATH,0.,1
EVENT. STOP,365.,1
PLOT. (PMEAN)
PLOT. (POP)
```

STATISTICAL REPORT FOR POP   TIME =   365.000000

TOTAL TIME	364.391079	AVERAGE	87.9511232	ST. DEV.	27.1806966
MAXIMUM	129.000000	MINIMUM	27.0000000	VARIANCE	738.790267
SUM	32048.6776	SUM SQ.	3087932.19	MEAN SQ.	8474.22553

STATISTICAL REPORT FOR BDEL   TIME =   365.000000

NUMBER	296.000000	AVERAGE	1.24070477	ST. DEV.	1.02023854
MAXIMUM	4.55162573	MINIMUM	.385895681	VARIANCE	1.04088668
SUM	367.248612	SUM SQ.	763.749563	MEAN SQ.	2.58023501

STATISTICAL REPORT FOR DDEL   TIME =   365.000000

NUMBER	369.00..L0	AVERAGE	.990517629	ST. DEV.	9.937122086E-02
MAXIMUM	1.28477325	MINIMUM	.693412332	VARIANCE	9.874639535E-03
SUM	365.501005	SUM SQ.	365.678931	MEAN SQ.	.990999813

PLOT NO. 1

PLOT NO. 2

Punching data decks.

The special utility subroutine PUNCHD provides the capability of obtaining a punched deck of the system-declared simulation control variables (refer to Table 2.1-1, with the exeception of FLOW) and all variables declared in STORAGE.. The data deck is punched in a format consistent for input to a SIMCOMP simulation (refer to section 2.2.1). Any variables which have not been assigned a value are ignored. The call to PUNCHD contains no parameters and is of the following form.

CALL PUNCHD

The call to PUNCHD can be placed at any point in the simulation, but will be executed only the first time it is called. The call in most cases should be placed in routines that are executed only once, such as START or FINIS. A call to routine PUNCHD is most useful when data is generated in one simulation and is to be used as input in another simulation. The punched deck produced by PUNCHD will have the same job number as the job that generated the punched deck.

2.1.6 Listing controls.

SIMCOMP normally produces a listing of the source section during compilation. The SIMCOMP compiler directives LIST. and NOLIST. may appear at any place in the source section. If a NOLIST. directive is encountered, the printing of all source statements from that point on is suppressed. The printing of source statements is reinitiated by the LIST. directive. LIST. and NOLIST. can appear anywhere in columns 1 through 72 and blanks are ignored.

2.1.7. Execution controls.

At CSU using the SCOPE 3.3 operating system, the
SIMCOMP compiler remains on the system as a permanent
file and is called up and executed by means of the job
control cards listed in Appendix C. SIMCOMP simulations
can be executed in any one of three different modes,
selected by the inclusion or absence of execution-
control directives in the source section. The three
modes of execution correspond to (i) the absence of
any execution directives, (ii) the inclusion of a DEBUG.
execution directive, and (iii) the inclusion of a NOGO.
execution directive. Either DEBUG. or NOGO. is key-
punched anywhere in columns 1 through 72 and blanks
are ignored. If both a DEBUG. and a NOGO. card appear
in the source section, then NOGO. is assumed.

The first mode of execution, that is, the default
action of the compiler, should be used during the early
stages of the development of a simulation. Any FORTRAN
compilation errors which are detected will be printed
in the output along with the offending statements. The
second mode of execution, selected by a DEBUG. directive,
should be used after all compilation errors have been
eliminated. If compilation errors are encountered while
DEBUG. has been selected, the only indication is a
message entered in the dayfile; the run is terminated.
The printed output will not contain any diagnostics

explaining the nature of the compilation error. The

use of DEBUG. is intended primarily for the detection

and reporting of execution errors. If DEBUG. is

selected and an arithmetic-mode error occurs during

the course of the simulation, a short explanation of

the error along with a dump of all variables and their

values is provided. A complete explanation of the

use of the DEBUG. facility is contained in section 2.3.

The third mode of execution, selected by the NOGO.

directive, suppresses compiler generation of the job

control cards and therefore requires the user to supply

the desired job control cards. When the default or

DEBUG. mode for execution is selected, a standard set

of job control cards are generated by the compiler

which will automatically execute the simulation. By

selecting NOGO. the user must supply his own control

cards if anything more than a SIMCOMP compilation is

desired. The job control card sequences generated in

the default and DEBUG. modes are listed in Appendix B.

2.1.8 Comments

Comments conform to the format for FORTRAN comments and can appear anywhere in the source section. Any statement with the letter C in column 1 will be taken as a comment. The commentary information can be any string of blanks and characters in columns 2 through 80. Each comment statement must begin with a C in column 1 and may not be continued by means of a non-blank character in column 6 on subsequent cards.

2.2 Data Section

When a SIMCOMP source program has been compiled
successfully, the first phase of execution of the simula-
tion begins by reading and processing the data section.
The data section is comprised of three types of state-
ments. These are (i) data value assignment statements,
(ii) output requests, and (iii) exogenous event requests.
All statements in the data section are free form in
columns 1 through 80 and blanks are ignored. Statements
within the data section can appear in any order. Illegally
formatted statements will produce a diagnostic, but in a
great majority of cases the errors are not fatal. An
attempt will be made to execute the simulation by assum-
ing default values for critical parameters. As a result
the output should be examined for data section diagnostics
if correct results are to be realized. The physical
location of the data section within the job deck is
illustrated in Appendix B.

2.2.1 Parameter input data.

The values of variables or arrays declared in
storage-allocation statements (refer to section 2.1.1)
and the values of the simulation-control variables
(refer to Table 2.1-1, with the exception of FLOW) may
be set by data-value assignment statements. Data-value
assignment is specified by statements of the following
form:

$$var = v \ \$$$

or

$$var = v_1, v_2, \ldots, v_n \ \$$$

The variable name or array element "var" must have been
declared in a storage-allocation statement or is a
reserved simulation-control variable. The value of the
variable or the values of the array "v" may be either
integer- or real-valued constants. The mode of the
value should correspond to the mode of the variable.
If the mode of the value and the variable differ the
mode of the value is converted to the mode of the vari-
able. If an attempt is made to assign a real value to
an integer variable, the value is truncated to an integer,
the assignment is made, and a diagnostic is issued.
Each expression is followed by a dollar sign. The
expressions are free form in columns 1 through 80 and
blanks are ignored. More than one expression may appear
on a single data card, in addition to being run on from

one card to the next. If any variable declared in a storage-allocation statement or state variable is not given a value in the data section, the variable is flagged as indefinite, and subsequent use of the variable prior to assigning the variable a value in the simulation will cause an arithmetic-mode error. Refer to section 2.3 on debugging for an explanation of the resulting diagnostic.

If the term "var" is an array element and a series of values are to be assigned to the array, the values are stored by columns in ascending order. This means that successive storage locations in a multiple-dimensioned array can be located by visualizing the left-most subscripts to vary the fastest. The array element "var" must be the location in the array where the storing of values begins. In an array declared as B(3,2,2) the order of the elements of the array is as follows:

$$
\begin{array}{l}
B(1,1,1) \\
B(2,1,1) \\
B(3,1,1) \\
B(1,2,1) \\
B(2,2,1) \\
B(3,2,1) \\
B(1,1,2) \\
B(2,1,2) \\
B(3,1,2) \\
B(1,2,2) \\
B(2,2,2) \\
B(3,2,2)
\end{array}
$$

For example, if we desired to set the last six locations of the above array to the values 1., 2., 3., 2., 4., and 6., respectively, we write:

$$B(1,1,2) = 1., 2., 3., 2., 4., 6. \$$$

If an array name is not followed by a subscript, the first element is assumed. That is,

$$B = 0., 1.22E2, 1.22E-2 \; \$$$

would result in

$$B(1,1,1) = 0.$$
$$B(2,1,1) = 1.22E2$$
and $$B(3,1,1) = 1.22E-2.$$

If a series of locations are to contain the same value, an integer-repetition factor may be used. If the last six elements of the array B were to have the values 1., 1., 1., 3., 3., and 9., respectively, then we write:

$$B(1,1,2) = 3*1., 2*3., 9. \; \$$$

The entire array could be set to zero with the single statement:

$$B = 12*0. \; \$$$

If a particular variable or array element is set more than once, the last assignment is assumed in effect. If the entire array above is to be set to zero with the exception of element B(2,2,1) which has the value 9.3×10^{-3}, we write:

$$B = 12*0. \; \$ \; B(2,2,1) = 9.3E-3 \; \$$$

Example 2.2.1-1. Illustration of data-value assignment.

```
STORAGE. S(3,3),S2(4),INDX
STORAGE. VA,VB,VC,TOP,PX,KVAL(2,3)
REAL. I(10),J(10)
INTEGER. T1,T2
```

78₉   end-of-record separator

```
S=9*1. $ VB=2. $ VC=3.14 $ VA=3.69385E-6 $ TOP=10.3 $ S(1,1)=0.2 $ S(2,2)=0.3 $
S(3,3)=0.5 $ INDX=1 $ KVAL=1,90,3*0,500 $ PX=9.9E+10 $ I=3*0,1 $ I(5)=5*2 $ J=1,
3,5,3*6,4*9 $ T1=20 $ T2=10 $ PX=999.999 $ S2(2)=39. $ S(4)=41. $ TSTRT=0. $ TEN
D=100. $ DT=0.01 $ DTPR=10. $ DTFL=20. $ X(1)=0. $ X(23)=3*50. $
```

<u>Simulation control variables</u>.

When a simulation has been defined containing flows, the reserved system control variables TSTRT, TEND, and DT should be given values in the data section. These variables must be given values in order for the simulation to execute; if the user fails to set the value of any or all of these three variables, default values will be assumed. Whenever the system assumes a default value for a control variable not specified by the user, a warning message is issued. Table 2.2.1-1 describes the default values assumed by the system. When the user supplies values for TSTRT, TEND, and DT the following conditions must hold:

$$TSTRT \leq TEND$$

$$DT > 0$$

If either or both of these conditions are not satisfied, a diagnostic is printed and the values in the last line of Table 2.2.1-1 are assumed.

Table 2.2.1-1. Default values of simulation control variables for simulations containing flows.

TSTRT	TEND	DT
given	given	(TEND-TSTRT)/10.
given	TSTRT + 10.*DT	given
TEND - 10.*DT	given	given
given	TSTRT + 10.	1.
TEND - 10.	given	1.
0.	10.*DT	given
0.	10.	1.

If the simulation does not contain any flows, the simulation is assumed to contain only events. In this case the reserved simulation control variable DT has no meaning and does not have to be given a value. In order for an event-only simulation to execute, the chain of events has to be initiated in either of two ways. If TSTRT is given a value in the data section and subroutine START has been supplied by the user, then execution time calls to the event schedule (refer to section 2.1.3) can be included in subroutine START to initiate the event sequence. If TSTRT is not given a value, then the only way an event sequence can be initiated is by including exogenous-event requests (refer to sections 2.1.3 and 2.2.3) in the data section. In this case if TSTRT has not been given a value, TSTRT is assumed to be equal to the time of the first exogenously scheduled event. If TEND is not given a value, no action is taken and the user has the responsibility of scheduling the system-defined event HALT at the appropriate time. If TEND is given a value, the simulation will be terminated automatically at time TEND. If TEND is less than or equal to the value of TSTRT, the value of TEND is assumed indefinite.

The reserved simulation control variables DTPR, DTPL, and DTFL are used by the SIMCOMP executive routine to determine the frequency of tabular print-outs, storage of values for plotting, and the printing of values

of flows respectively. If one of the above output actions

is requested (refer to section 2.2.2) and the correspond-

ing simulation control variable has not been given a

value in the data section or was given an illegal value,

the system will issue a diagnostic and a default value

will be chosen. The following default values (Table

2.2.1-2) are those chosen if at least one flow has been

defined in the simulation.

Table 2.2.1-2. Default values of output control variables for continuous
                simulations (containing flow definitions).

Condition	Default Action
PRINT. request(s) present and	
(1) $DTPR \leq 0.$ or indefinite.	$DTPR$ = maximum of (TEND-TSTRT)/10 and DT.
(2) $DTPR < DT.$	$DTPR = DT.$
PLOT. request(s) present and	
(1) $DTPL \leq 0.$ or indefinite; and TIME is the independent vari- able of a least one plot.	$DTPL$ is set to a value which provides good readability while minimizing execution time.
(2) $DTPL \leq 0.$ or indefinite; and TIME is not the independent vari- able of any plot.	$DTPL = DT.$
(3) $DTPL < DT.$	$DTPL = DT.$
FLOW. request(s) present and	
(1) $DTFL \leq 0.$ or indefinite.	$DTFL$ = maximum of (TEND-TSTRT)/10 and DT.
(2) $DTFL < DT.$	$DTFL = DT.$

If no flows have been defined in the simulation and only events have been defined, the following table (Table 2.2.1-3) describes the default values chosen for the output simulation control variables.

Table 2.2.1-3. Default values of output control variables for simulations containing only events.

Condition	Default Action
PRINT. request(s) present and	
(1) DTPR $\leq$ 0 or indefinite, and TEND is defined.	DTPR = (TEND-TSTRT)/10.
(2) DTPR $\leq$ 0 or indefinite, and TEND is undefined.	DTPR = 1.
PLOT. request(s) present and	
(1) DTPL $\leq$ 0 or indefinite; TEND is defined; and TIME is the independent variable of a least one plot.	DTPL is se to a value which provides good readability while minimizing execution time.
(2) DTPL $\leq$ 0 or indefinite; TEND is defined; and TIME is not the indefinite variable of any plot.	DTPL = 1.
(3) TEND is undefined.	DTPL = 1.

2.2.2 Output requests.

The results of a simulation can be requested as printed tables of values through time in addition to printed or microfilm plots. Output requests also allow a measure of control over the printing of the initial values of variables declared in STORAGE.. An execution trace facility is also provided which is especially useful during the debugging stages of simulations containing events. The format and usage of these commands are described in the following pages. In general, an output request is comprised of a request verb followed by a period with the remainder of the card containing the necessary information. The output requests are free form in columns 1 through 80 with blanks ignored. All output requests are contained in the data section.

PRINT. requests.

Printed tabular output requests are specified by the statement form:

$$\text{PRINT. } v_1, v_2, \ldots, v_m$$

Each of the variables v_i to be printed must be a variable or array location which appears in a storage allocation statement (i.e., STORAGE., INTEGER., or REAL. statement) in the source section. The values of the state variables may also be requested for printing by specifying "X(i)" where i is the compartment number.

If a card with the word PRINT. without a list of vari-

ables is included, then all state variables will be

printed. As many PRINT. cards as needed may be included.

The variable which controls the frequency of

tabular output through simulated time is DTPR. The

value for DTPR should be set by means of a data-assign-

ment expression. If DTPR is not given a value, a

diagnostic will be issued and a default value assumed.

Tables 2.2.1-2 and 2.2.1-3 describe the values chosen.

The following example illustrates some legal PRINT.

requests. Presumably the state variables requested in

the second request would have been defined in the

simulation. An example of the output produced by PRINT.

requests is presented in example 2.2.2-1.

Example 2.2.2-1. Illustration of PRINT. requests.

```
STORAGE. RSET,QVAL(3),TQ(4,3,2)
REAL. NPOP,IST(2,3)
INTEGER. DAY,MONTH,YEAR
```
7_89 *end-of-record separator*
```
PRINT. DAY,MONTH,YEAR,RSET,IST(1,1),IST(2,1),QVAL(1),QVAL(2),QVAL(3),TQ(1,1,1)
PRINT. NPOP,X(3),X(4),X(66),X(976)
PRINT.
```

FLOW. requests.

The printing of computed values for the flows

over each integration step is specified by the statement

form:

FLOW. (i,j), (k,l), ... , (m,n)

Each parenthesized pair of numbers refers to a
flow defined in the source section. If a card with
the word FLOW. is included without a list of particular
flows, then all flows defined in the simulation will
be printed. If a requested flow (i,j) has not been
defined, a diagnostic is issued and the illegal request
is ignored.

The variable which controls the frequency with
which flows are printed is DTFL. The value for DTFL
should be set by means of a data assignment expression.
If DTFL is not given a value, a diagnostic will be
issued and a default value assumed. Tables 2.2.1-2 and
2.2.1-3 describe the values chosen. When evaluating
the performance of a simulation, care must be taken to
associate the values of the flows with the correct time
step. The following simple example illustrates the
relationship between the times at which the state vari-
ables are printed via a PRINT. request and the times
at which the flows are printed via a FLOW. request.

Example 2.2.2-2. Illustration of FLOW. requests.

```
(100-200). FLOW=0.01*X(200)
(I=1,3-500). FI=I*2
      FLOW=0.1*X(I)/FI
```

$7 8_9$ *end-of-record separator*

```
X(1)=3*100. $ X(100)=1000. $ X(200)=1. $ X(500)=0. $
TSTRT=0. $ TEND=10. $ DT=1. $ DTFL=1. $ DTPR=1. $
PRINT.
FLOW.
```

SIMULATION RESULTS
(partial listing)

TIME = 0.
        X(100) = 1000.00000   X(200) = 1.00000000   X(1) = 100.000000   X(500) =   0
        X(2) = 100.000000   X(3) = 100.000000

VALUES OF FLOWS. TIME = 0.   TO 1.00000000
        FLOW(100.200) = .10000E-01 FLOW(1.500) = 5.00000000 FLOW(2.500) = 2.50000000 FLOW(3.500) = 1.66666667

TIME = 1.00000000
        X(100) = 999.990000   X(200) = 1.01000000   X(1) = 95.0000000   X(500) = 9.16666667
        X(2) = 97.5000000   X(3) = 98.3333333

VALUES OF FLOWS. TIME = 1.00000000   TO 2.00000000
        FLOW(100.200) = .10100E-01 FLOW(1.500) = 4.75000000 FLOW(2.500) = 2.43750000 FLOW(3.500) = 1.63888889

TIME = 2.00000000
        X(100) = 999.979900   X(200) = 1.02010000   X(1) = 90.2500000   X(500) = 17.9930556
        X(2) = 95.0625000   X(3) = 96.6944444

VALUES OF FLOWS. TIME = 2.00000000   TO 3.00000000
        FLOW(100.200) = .10201E-01 FLOW(1.500) = 4.51250000 FLOW(2.500) = 2.37656250 FLOW(3.500) = 1.61157407

TIME = 3.00000000
        X(100) = 999.969699   X(200) = 1.03030100   X(1) = 85.7375000   X(500) = 26.4936921
        X(2) = 92.6859375   X(3) = 95.0828704

VALUES OF FLOWS. TIME = 3.00000000   TO 4.00000000
        FLOW(100.200) = .10303E-01 FLOW(1.500) = 4.28687500 FLOW(2.500) = 2.31714844 FLOW(3.500) = 1.58471451

TIME = 4.00000000
        X(100) = 999.959396   X(200) = 1.04060401   X(1) = 81.4506250   X(500) = 34.6824301
        X(2) = 90.3687891   X(3) = 93.4981559

VALUES OF FLOWS. TIME = 4.00000000   TO 5.00000000
        FLOW(100.200) = .10406E-01 FLOW(1.500) = 4.07253125 FLOW(2.500) = 2.25921973 FLOW(3.500) = 1.55830260

TIME = 5.00000000
        X(100) = 999.948990   X(200) = 1.05101005   X(1) = 77.3780937   X(500) = 42.5724836
        X(2) = 88.1095693   X(3) = 91.9398533

VALUES OF FLOWS. TIME = 5.00000000   TO 6.00000000
        FLOW(100.200) = .10510E-01 FLOW(1.500) = 3.86890469 FLOW(2.500) = 2.20273923 FLOW(3.500) = 1.53233089

TIME = 6.00000000
        X(100) = 999.938480   X(200) = 1.06152015   X(1) = 73.5091891   X(500) = 50.1764585
        X(2) = 85.9068301   X(3) = 90.4075224

VALUES OF FLOWS. TIME = 6.00000000   TO 7.00000000
        FLOW(100.200) = .10615E-01 FLOW(1.500) = 3.67545945 FLOW(2.500) = 2.14767075 FLOW(3.500) = 1.50679204

TIME = 7.00000000
        X(100) = 999.927865   X(200) = 1.07213535   X(1) = 69.8337296   X(500) = 57.5063807
        X(2) = 83.7591593   X(3) = 88.9007303

VALUES OF FLOWS. TIME = 7.00000000   TO 8.00000000
        FLOW(100.200) = .10721E-01 FLOW(1.500) = 3.49168648 FLOW(2.500) = 2.09397898 FLOW(3.500) = 1.48167884

TIME = 8.00000000
        X(100) = 999.917143   X(200) = 1.08285671   X(1) = 66.3420431   X(500) = 64.5737250
        X(2) = 81.6651804   X(3) = 87.4190515

VALUES OF FLOWS. TIME = 8.00000000   TO 9.00000000
        FLOW(100.200) = .10828E-01 FLOW(1.500) = 3.31710216 FLOW(2.500) = 2.04162951 FLOW(3.500) = 1.45698419

PLOT. requests.

> Plotted output requests are specified by the
> following statement forms:
>
> $$\text{PLOT. } (\text{group}_1), \ldots, (\text{group}_n)$$
>
> or
>
> $$\text{PLOT. } (\text{group}_1), \ldots, (\text{group}_n) \text{ phrase}$$

Each plot card will generate a single plot (on

the line printer unless a FILM. card is included--refer

to page 2.2.2-12) of the variables listed in the groups

on the plot card. Each "(group$_i$)" is an expression

of the following form:

$$(u_1, \ldots, u_m)$$

or

$$(u_1, \ldots, u_m [min,max])\underline{5/}$$

Each of the terms "u$_i$" is an expression specifying

a dependent variable and takes on one of the following

forms:

var

or

var=c

or

var.LOG

or

var.LOG=c

Each term "var" is a dependent variable to be

plotted and must be a simple variable or location in

an array which was declared in a storage-allocation

statement in the source section, or is a state variable

of the form "X(i)" where i is a compartment number.

The logorithm to the base 10 of a variable is plotted

by specifying ".LOG" immediately following the variable

$\underline{5/}$ The characters "[" and "]" are represented on a key punch by the
multipunches 8.7 card 0-8-2, respectively.

name and its subscripts if present. In this case the
variable is plotted on a log scale, but the true values
of the variable are printed on the dependent axis. If
one variable in a group is requested to be plotted on
a log scale, all variables in the group will be so plot-
ted. The character used in the plot to identify the
particular variable is normally chosen by the plotting
routine. An index of the variables plotted and the
characters used is printed out. A specific character
for a variable can be selected by appending the expression
"=c," where "c" is the character to be plotted.

Each group of variables (group$_i$) is scaled on
the plot independently of the other groups. One through
five groups per plot card with one through five variables
per group are allowed. If the expression in brackets
"[min,max]" is included in a group, the minimum and
maximum specified in the brackets are used as the extremes
of the dependent variable(s) in the group. The terms
"min" and "max" must be integer- or real-valued constants.
If the extreme values for a group are not specified,
the minimum and maximum values for all variables in the
group, appropriately rounded for readability, are used
as the extreme values of the group.

The optional modifying expression "phrase" is of
the following forms:

/ var

or

/ var [min,max]

or

[min,max]

The term "/ var" specifies the independent variable
for the plot. Only one independent variable is allowed
per plot. If the term "/ var" is omitted, "/ TIME" is
assumed. The term "var" can be any variable declared
in a storage-allocation statement or a state variable of
the form $X(i)$. The bracketed expression "[min,max]"
specifies the extreme values to be used in the plot
for the independent variable. If the bracketed expres-
sion is omitted, the minimum and maximum values for
the independent variable are used. The quantities "min"
and "max" must be integer- or real-values constants.

When PLOT. cards have been included in the data
section, the values of the variables named on PLOT.
cards are saved on a mass storage device during the
simulation at intervals of DTPL. These values are later
retrieved and used to produce the plots. As long as
TIME is the independent variable for at least one plot,
the value of DTPL does not have to be set by the user
via a data-assignment statement. A value for DTPL
which reduces the amount of data to be saved while
preserving the ultimate readability of the plot is
chosen automatically. Refer to Tables 2.2.1-2 and
2.2.1-3 for a complete description of the choices of
values for DTPL.

PLOT. requests are free form in columns 1 through
80 and blanks are ignored. Each plot request must be
completed on one data card. The following examples
illustrate some of the possibilities of formatting plot
requests. All of the printer plots were generated in
the same run and are reproduced following the examples
of the plot requests.

Example 2.2.2-3.   Illustration of PLOT. requests with printer-plotted output.

```
STORAGE. T,A,A1,RX,RY,TSIN
STORAGE. U,V
(1-2).
        T=TIME*6.28/50.
        RX=A*T-A1*SIN(T)
        RY=A-A1*COS(T)
        TSIN=SIN(T)
        FLOW=?.*TSIN
        V=TIME/10.
        U=EXP(V*SIN(V))
```

7_89 *end-of-record separator*

```
TSTRT=0. $ TEND=100. $ DT=1. $ X=2*0. $
T=0. $ A=1. $ A1=2. $ RX=0. $ RY=-1. $ TSIN=0. $
V=0. $ U=1. $
PLOT. (RX=X,RY=Y)/T
PLOT. (RY=*[-1.5,3.5])/RX[-2,14.5]
PLOT. (TSIN).(X(1))
PLOT. (TSIN=*[-2,2])/X(1)[-40,+40]
PLOT. (U)/V
PLOT. (U.LOG)/V
```

GRAPHICAL SIMULATION RESULTS   07/23/73   18.48.56.

GRAPH NO.	GROUP	GROUP RANGE DECLARATION	DEPENDENT VARIABLE(S)	PLOTTED CHARACTER	INDEPENDENT VARIABLE	INDEPENDENT VARIABLE RANGE DECLARATION
1	1		RX RY	X Y	T	
2	1	-1.50 TO 3.50	RY	*	RX	-2.00 TO 14.5
3	1 2		TSIN X(1)	A B	TIME	
4	1	-2.00 TO 2.00	TSIN	*	X(1)	-40.0 TO 40.0
5	1		U	A	V	
6	1		U	LOG A	V	

PLOT NO. 1

FLOT NO. 2

PLOT NO. 3

PLOT NO. 4

PLOT NO. 5

PLOT NO. 6

FILM. requests.

PLOT. requests will cause the graphs to be
produced on the line printer and will accompany the
output. A FILM. card included in the data section,
all plots requested will be generated on microfilm.
Plots generated on microfilm will generally have a
higher degree of resolution than those produced by the
line printer. The FILM. card is free form in columns
1 through 80 and blanks are ignored. The plots which
were illustrated in the preceding section are reproduced
in the following from microfilm.

Example 2.2.2-4. A sample of microfilm output.

GRAPHICAL SIMULATION RESULTS   07/23/73   18.50.54.

GRAPH NO.	GROUP	GROUP RANGE DECLARATION	DEPENDENT VARIABLE(S)	PLOTTED CHARACTER	INDEPENDENT VARIABLE	INDEPENDENT VARIABLE RANGE DECLARATION
1	1		RX RY	X Y	Y	
2	1	-1.50 TO 3.50	RY	•	RX	-2.00 TO 4.5
3	1		TSIN	A	TIME	
	2		X(1)	B		
4	1	-2.00 TO 2.00	TSIN	•	X(1)	-40.0 TO 40.0
5	1		U	A	V	
6	1		U	LOG A	V	

PLOT NO. 1

PLOT NO. 2

PLOT NO. 3

PLOT NO. 4

PLOT NO. 5

PLOT NO. 6

TITLE. requests.

A one-line title for each plot can be requested
by a statement of the following form:

TITLE.   text

A TITLE. card is free format in columns 1 through
80. The title for each plot is indicated by placing a
TITLE. card before the corresponding PLOT. card. If a
TITLE. card appears after all PLOT. requests or if more
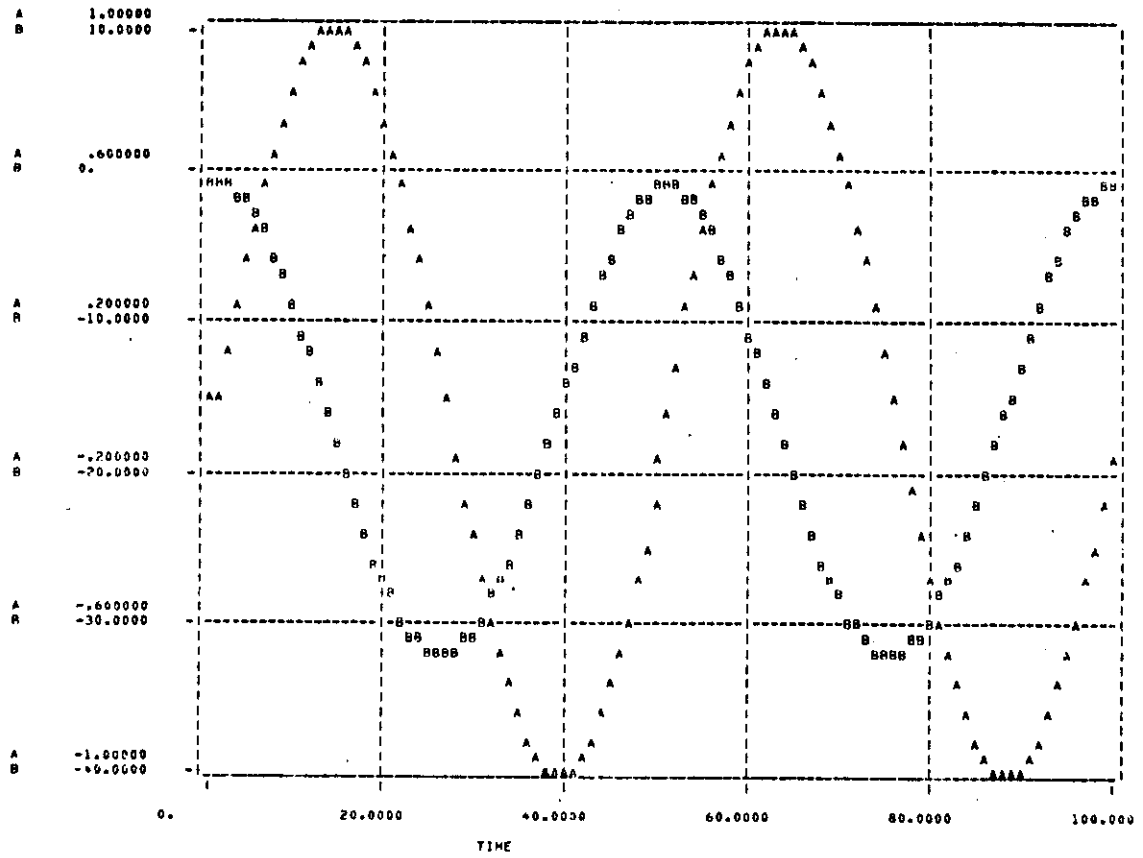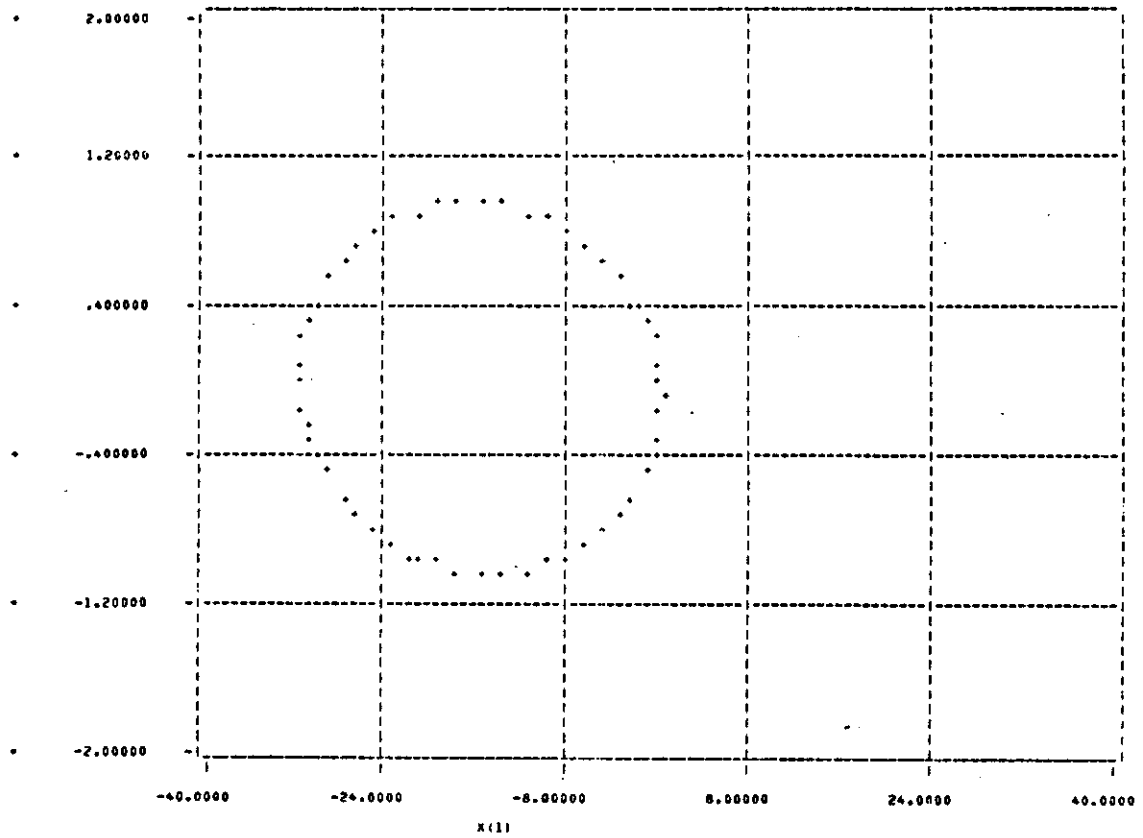than one TITLE. card precedes a PLOT. request, a diag-
nostic is issued and the offending TITLE. card or cards
will be ignored. The TITLE. card immediately preceding
a PLOT. card will be used as the title of the plot.
Any nonblank characters following the period on a title
card is assumed part of the title and will be reproduced
at the top of the corresponding plot. The use of TITLE.
requests is illustrated in example 2.1.3-1.

Initial-conditions listing controls.

After the data section has been processed but
prior to the start of execution of the simulation,
the initial values of the simulation control variables,
state variables, and primary user-declared variables
are printed. Primary user-declared variables are those
variables declared in storage-allocation statements
which are not prefaced by an asterisk (refer to section
2.1.1). The initial values of secondary user-declared
variables (i.e., those variables prefaced by an asterisk

in storage-allocation statements), in addition to the

initial values normally printed, are requested by the

inclusion of an ALL. card in the data section. The

characters ALL. are free form in columns 1 through

80, and the period must be included. Similarly, the

printing of *all* initial conditions is suppressed by

the inclusion of a NONE. card. The characters NONE.

are free form in columns 1 through 80, and the

trailing period must be included.

Both of the examples presented are followed by the

initial-conditions output which was requested.

Example 2.2.2-5. An example of normal initial-conditions output.

```
STORAGE. S(3,3),S2(4),INDX,*LARGE(25)
STORAGE. VA,VB,VC,TOP,PX,KVAL(2,3)
REAL. I(10),J(10)
INTEGER. T1,T2,*TEMP1,*TEMP2
```
78₉   *end-of-record separator*
```
S=9*1. $ VB=2. $ VC=3.14 $ VA=3.69385E-6 $ TOP=10.3 $ S(1,1)=0.2 $ S(2,2)=0.3 $
S(3,3)=0.5 $ INDX=1 $ KVAL=1,90,3*0,500 $ PX=9.9E+10 $ I=3*0,1 $ I(5)=5*2 $ J=1,
3,5,3*6,4*9 $ T1=20 $ T2=10 $ PX=999.999 $ S2(2)=39. $ S(4)=41. $ TSTRT=0. $ TEN
D=100. $ DT=0.01 $ DTPR=10. $ DTFL=20. $ X(1)=0. $ X(23)=3*50. $
LARGE=2,85,63,45,72,15,38,92,100,12,3,4,50,5*10 $ LARGE(21)=5*1 $
```

SIMCOMP VERSION 3.0   PARAMETER VALUES

- SIMULATION CONTROL PARAMETERS -

```
TSTRT =           0
TEND  = 100.000000
DT    = .100000000E-01
DTPR  = 10.0000000
DTFL  = 20.0000000
```

- PRIMARY USER DEFINED VARIABLES -

```
I(1-3) =          0      I(4) = 1.00000000      I(5-9) = 2.00000000      I(10) = INDEFINITE
INDX =            1      J(1) = 1.00000000      J(2) = 3.00000000        J(3) = 5.00000000
J(4-6) = 6.00000000     J(7-10) = 9.00000000   KVAL(1,1) =          1    KVAL(2,1) =        90
KVAL(1,2-1,3) =    0.   KVAL(2,3) =       500   PX = 999.999000          S(1,1) = .200000000
S(2,1-3,1) = 1.00000000 S(1,2) = 41.0000000    S(2,2) = .300000000      S(3,2-2,3) = 1.00000000
S(3,3) = .500000000     S2(1) = INDEFINITE      S2(2) = 39.0000000       S2(3-4) = INDEFINITE
TOP = 10.3000000        T1 =            20      T2 =           10        VA = .369385000E-05
VB = 2.00000000         VC = 3.14000000
```

Example 2.2.2-6. An example of initial-conditions output using ALL.

```
STORAGE.  S(3,3),S2(4),INDX,*LARGE(25)
STORAGE.  VA,VB,VC,TOP,PX,KVAL(2,3)
REAL.  I(10),J(10)
INTEGER.  T1,T2,*TEMP1,*TEMP2
```

7_{8_9} *end-of-record separator*

```
S=9*1.  $  VB=2.  $  VC=3.14  $  VA=3.69385E-6  $  TOP=10.3  $  S(1,1)=0.2  $  S(2,2)=0.3  $
S(3,3)=0.5  $  INDX=1  $  KVAL=1,90,3*0,500  $  PX=9.9E+10  $  I=3*0,1  $  I(5)=5*2  $  J=1,
3,5,3*6,4*9  $  T1=20  $  T2=10  $  PX=999.999  $  S2(2)=39.  $  S(4)=41.  $  TSTRT=0.  $  TEN
D=100.  $  DT=0.01  $  DTPR=10.  $  DTFL=20.  $  X(1)=0.  $  X(23)=3*50.  $
LARGE=2,85,63,45,72,15,38,92,100,12,3,4,50,5*10  $  LARGE(21)=5*1  $
ALL.
```

SIMCOMP VERSION 3.0   PARAMETER VALUES

          - SIMULATION CONTROL PARAMETERS -

               TSTRT =   0
               TEND = 100.000000
               DT   = .100000000E-01
               DTPR = 10.0000000
               DTFL = 20.0000000

          - PRIMARY USER DEFINED VARIABLES -

I(1-3) = 0	I(4) = 1.00000000	I(5-9) = 2.00000000	I(10) = INDEFINITE
INDX = 1	J(1) = 1.00000000	J(2) = 3.00000000	J(3) = 5.00000000
J(4-6) = 6.00000000	J(7-10) = 9.00000000	KVAL(1,1) = 1	KVAL(2,1) = 90
KVAL(1,2-1,3) =	KVAL(2,3) = 500	PX = 999.999000	S(1,1) = .200000000
S(2,1-3,1) = 1.00000000	S(1,2) = 41.0000000	S(2,2) = .300000000	S(3,2-2,3) = 1.00000000
S(3,3) = .500000000	S2(1) = INDEFINITE	S2(2) = 39.0000000	S2(3-4) = INDEFINITE
TOP = 10.3000000	T1 = 20	T2 = 10	VA = .369385000E-05
VB = 2.00000000	VC = 3.14000000		

          - SECONDARY USER DEFINED VARIABLES -

LARGE(1) = 2	LARGE(2) = 85	LARGE(3) = 63	LARGE(4) = 45
LARGE(5) = 72	LARGE(6) = 15	LARGE(7) = 38	LARGE(8) = 92
LARGE(9) = 100	LARGE(10) = 12	LARGE(11) = 3	LARGE(12) = 4
LARGE(13) = 50	LARGE(14-18) = 10	LARGE(19-2J) = INDEFINITE	LARGE(21-25) = 1
TEMP1 = INDEFINITE	TEMP2 = INDEFINITE		

Event execution trace.

Simulations containing a large number of events,
in which the logical structure for scheduling and
rescheduling the events is complicated, are sometimes
difficult to debug. It is important in such cases to
determine that the events are being scheduled and
executed in the proper sequence. The event-execution

trace facility is provided to aid in debugging this
type of simulation. The trace facility is envoked
by including in the data section a card of the following
form:

TRACE.

The command is free form in columns 1 through 80.
Blanks are ignored.

The trace facility will print in the output the
contents of the event stack, including event names,
scheduled times of occurrence, and priorities, each time
an event is executed. The current value of simulated
time is also printed. Care should be taken in using
the trace feature since in some simulations a very
large amount of output can be produced.

2.2.3. Exogenous event requests.

Events can be scheduled externally prior to the start of simulation by statements of the following form:

EVENT. name, time, priority

The event name must be the name of an event defined in the source section or one of the system-defined events included in Table 2.1.3-1. The event HALT is the most commonly used system-defined event. The simulated time of occurrence, "time," must be either an integer- or real-valued constant. Integer-valued constants are converted to real-valued constants internally since TIME is a real-valued variable. The priority of the event, "priority," must be an integer- or real-valued constant in the range 1 to 512. Real-valued priorities specified on an event card are truncated to an integer. Priorities outside the range 1 through 512 are assumed to be 512. If the priority is not specified, a priority of one is assumed.

A maximum of 20 exogenous event requests can be included in the data section. If more than 20 requests are needed, endogenous event requests of any number could be included in subroutine START (refer to section 2.1.3). An example of an exogenous event request is contained in example 2.1.1-2.

2.3 Debugging.

During the course of programming a simulation, three phases of development occur. These are (i) the detection and correction of compilation errors, (ii) the detection and correction of execution errors, and (iii) the evaluation and refinement of the results of the simulation in preparation for production runs. Section 2.1.7 described the use of special execution controls. The default mode of execution should be used during the first phase of development. Once compilation errors have been eliminated, the DEBUG. mode of execution should be selected. The inclusion of a DEBUG. statement in the source section enables the simulation executive routine to detect arithmetic mode errors and produce a report. Arithmetic mode errors occur when illegal values are used in an arithmetic operation. Other types of errors can occur which will not produce a report, but are usually self explanatory such as an exceeded time limit.

Types of errors reported.

As stated above, arithmetic mode errors are detected when an illegal value is used in a computation. Table 2.3-1 can be the result of an operation, but will not be detected and reported until the resulting illegal value is used as an operand in a computation. Table 2.3-1 summarizes the illegal conditions which are detected and reported.

Table 2.3-1. Summary of error modes.

Error Mode	Condition
1	Address out of range - an attempt was made to reference central memory outside of established limits.
2	Operand out of range - the floating point arithmetic unit attempted to use an infinite operand.
3	Combined errors 1 and 2.
4	Indefinite operand - floating point arithmetic unit attempted to use an indefinite operand.
5	Combined errors 1 and 4.
6	Combined errors 2 and 4.
7	Combined errors 1, 2, and 4.

Error number one, address out of range, usually occurs when an index of a subscripted variable gets too large. Error numbers two and four occur when infinite or indefinite operands are used in the computation. The possible ways in which infinite or indefinite operands can be generated by division are illustrated in Table 2.3-2, using the following definitions of floating point (real) values. In Table 2.3-3, "X" represents any octal digit.

Table 2.3-2. Illegal results produced by division (A/B).

		B			
		+N	-N	+0	-0
A	+N	-	-	$+\infty$	$-\infty$
	-N	-	-	$-\infty$	$+\infty$
	+0	0	0	+IND	+IND
	-0	0	0	+IND	+IND

Table 2.3-3. Definition of floating point operands and results.

Mnemonic	Octal(internal) Representation	Meaning
+0	0000 X ... X	positive zero
-0	7777 X ... X	negative zero
$+\infty$	3777 X ... X	positive infinite
$-\infty$	4000 X ... X	negative infinite
+IND	1777 X ... X	positive indefinite
-IND	6000 X ... X	negative indefinite
N	--	any value with the exception of $\pm\infty$, $\pm$IND, or $\pm$0.

Positive or negative infinite results can be generated whenever a computation yields a floating point value whose absolute magnitude is outside the range 10^{-293} to 10^{322}. Such a condition can occur in

iterative computations where the value of a variable grows exponentially. This case can happen quite easily in simulations where a flow is defined to be proportional to a state variable which is linked by the flow.

Variables which are declared in storage-allocation statements (i.e., STORAGE., INTEGER., and REAL. statements) or reserved simulation control variables including state variables which are used in the computation, but have not been given a value will cause a mode 4 error. All such variables are initialized to an indefinite value before the data section is processed. Assigning a variable a value in the data section or in the simulation prior to the use of the variable as an operand will avoid the detection of an indefinite value.

An indefinite operand will also be detected if the variable FLOW is not assigned a legal value within the range of each flow declaration. The source and destination state variables which are linked by flows must be assigned legal values prior to the start of simulation, or an indefinite operand will be detected.

Debug reports.

The following sample simulation is shown to illustrate the information contained in a debug report. The listing in example 2.3-1 is the source and data section used which produced the debug report which follows.

Example 2.3-1.   A sample simulation containing an error.   Note: circled

numbers refer to items explained in the text.

```
DEBUG.
STORAGE. P.Q.R
(10-12). CALL PVAL
       FLOW=P*COS(TIME*6.28/50.)
       SUBROUTINE PVAL
       P=TIME/R+Q
       RFTURN
       END
```

7_89 *end-of-record separator*

TSTRT=0. \$ TEND=100. \$ DT=1. \$ X(10)=100. \$ X(12)=0. \$ R=2. \$ DTPR=20. \$
PRINT.

ARITHMETIC MODE ERROR   DIAGNOSTIC DUMP   07/23/73   18.55.59.

TYPE OF ERROR: ①
         ERROR MODE = 4   ②

①   FLOATING POINT ARITHMETIC UNIT RECIEVED AN INDEFINITE OPERAND

         OCCURING (APPROXIMATELY) AT ADDRESS 0076228 WHICH IS LOCATION 000004B IN ROUTINE PVAL ③

         NON-STANDARD FLOATING POINT ARITHMETIC - TABLES OF NON-STANDARD RESULTS BY DIVISION

              DIVIDE (A/B)   WHERE

                      B   +0   =   0000 X...X B
                 +N   -N   +0   -0   -0   =   7777 X...X B
                                                    +INF =   3777 X...X B   (+ INFINITY)
            +N   --   -- +INF -INF   -INF =   4000 X...X B   (- INFINITY)
          A -N   --   -- -INF +INF   +IND =   1777 X...X B   (+ INDEFINITE)
            +0   0   0   +IND +IND   -IND =   6000 X...X B   (- INDEFINITE)
            -0   0   0   +IND +IND   N   =   ANY WORD EXCEPT +INF, -INF, +IND, -IND, +0, OR -0

EXCHANGE JUMP PACKAGE: ⑤   ⑥   ⑦   ⑧   ⑨   ⑩

  ADDRESS   REFERENCED   TYPE   INDEX LOCAL   CONTAINED   DECODED VALUE
  REGISTERS CONTENTS VARIABLE NAME **ARRAY (DEC) ADDRESS   IN   VALUE (OCTAL)

                                                                      OUT OF RANGE
   A0   052000B   0000158 XFLOWS   0000 0000 0000 0000 0000B   0
   A1   0076108   000101B XEXECTV 0000 0000 0000 0001 0663B   4531
   A2   006136B   000010B XTRNSF   0100 0000 0046 0004 60008   1.3084312383453-280
   A3 ④ 010422B   R   REAL   0017618 / /   1721 4000 0000 0000 00008   2.000000000000
   A4   016717R   Q   REAL   001760B / /   1777 0000 0000 0000 1761B   + INDEFINITE
   A5   016716R   FLOW   REAL   0000228 XFLOWS   1777 0000 0000 0000 0001B   + INDEFINITE
   A6   0076158   XMFL   INTEGER   0000218 XFLOWS   0000 0000 0000 0000 00008   •
   A7   0076148

⑪   OPERAND   INCREMENT
    REGISTERS   CONTENTS   DECODED VALUE   REGISTERS   CONTENTS

     X0   0000 0000 0000 0000 0000R   0   B0   0000008
     X1   0000 0000 0000 0000 0000R   0   B1   0115308
     X2   0000 0106 6300 0000 00008   486512420454   B2   0000135
     X3 ⑫ 0100 0000 0046 0004 60008   1.3084312383453-280 ⑫   B3 ⑬   0000138
     X4   1721 4000 0000 0000 0000R   2.000000000000   B4   0000008
     X5   1777 0000 0000 0000 17618   + INDEFINITE   B5   0000018
     X6   1777 0000 0000 0000 0001B   + INDEFINITE   B6   0000018
     X7   1777 0000 0000 0000 0000R   + INDEFINITE   B7   0115278
```

VARIABLE DUMP  -  XFLOWS

| VARIABLE NAME | TYPE *=ARRAY | LOCATION | LOCAL ADDRESS | REPEATED | VALUE (OCTAL) | DECODED VALUE |
|---|---|---|---|---|---|---|
| XMFL | INTEGER | 0076148 | 0000218 | | 0000 0000 0000 0000 0000B | 0 |
| FLOW | REAL | 0076158 | 0000228 | | 1777 0000 0000 0000 0000B | + INDEFINITE |

VARIABLE DUMP  -  / /

 (III)

| VARIABLE NAME | TYPE *=ARRAY | LOCATION | LOCAL ADDRESS | REPEATED | VALUE (OCTAL) | DECODED VALUE |
|---|---|---|---|---|---|---|
| XADRS | *REAL | 0147368 | 000000R | | 5663 0663 3606 5301 3470B | -1.9006556451888E+37 |
| TIME | REAL | 0147378 | 0000018 | | 0000 0000 0000 0000 0000B | 0. |
| TSTRT | REAL | 0147408 | 0000028 | | 0000 0000 0000 0000 0000B | 0. |
| TEND | REAL | 0147418 | 0000038 | | 1726 6200 0000 0000 0000B | 100.0000000000 |
| DT | REAL | 0147428 | 000004R | (14) | 1720 4000 0000 0000 0000B | 1.000000000000 |
| DTPR | REAL | 0147438 | 000005R | | 1724 5000 0000 0000 0000H | 20.00000000000 |
| DTPL | REAL | 0147448 | 0000068 | | 1777 0000 0000 0000 0000B | + INDEFINITE |
| DTFL | REAL | 0147458 | 0000078 | | 1777 0000 0000 0000 0000B | + INDEFINITE |
| X | *REAL | 0147468 | 0000108 | 9 _ | 1777 0000 0000 0000 0000B | + INDEFINITE |
| | | | | | 1726 6200 0000 0000 0000B | 100.0000000000 |
| | | | | | 1777 0000 0000 0000 0000B | + INDEFINITE |
| | | | | | 0000 0000 0000 0000 0000B | 0. |
| | | | | 987 _ | 1777 0000 0000 0000 0000R | + INDEFINITE |
| P | REAL | 0167158 | 0017578 | | 1777 0000 0000 0000 0000B | + INDEFINITE |
| Q | REAL | 0167168 | 0017608 | | 1777 0000 0000 0000 0000B | + INDEFINITE |
| R | REAL | 0167178 | 0017618 | | 1721 4000 0000 0000 0000B | 2.000000000000 |
| XEVSTK | *REAL | 0167208 | 0017628 | | 0074 0000 0010 0001 0424B | 1.7216235906173-282 |

All debug reports contain the following items which refer to the circled numbers in example 2.3-1. A debug report contains three parts. These are (I) an explanation of the type of error, (II) the exchange jump package, and (III) the values of all variables in the simulation when the error was detected.

The first part contains (1) the error mode and (2) an explanation of the error mode (refer to table 2.3-1). After this information, (3) the routine in which the error was detected is listed. In this example an attempt was made to use an indefinite quantity in subroutine PVAL. We can immediately infer that either an indefinite value was generated by an undefined operation earlier in the simulation (such as 0 ÷ 0), or an attempt was made to use a variable which was never initialized. In this example the error was detected in the user-supplied routine PVAL. It is possible to

have errors detected in the simulation executive routines.
If an error is detected in routine XFLOWS, the illegal
condition was detected while a flow was being computed.
The executive routine which updates the state variables
is called XCSIM. An error in this routine indicates
that either the value of a flow is illegal or the value
of a state variable is illegal. Errors can also be
detected within FORTRAN-intrinsic functions such as ALOG
and EXP and will be reported accordingly.

The second part of a debug report is the exchange
jump package. This portion of the report reflects
the contents and meaning of the operation registers
at the time the error was detected. All computations
in the computer are accomplished by operating on
the values in these registers. The address registers
(4) contain the addresses in central memory of vari-
ables whose values were currently being used or were
recently used. If an address corresponds to the
address of a user variable, the name (5), the mode (6),
the one-dimensional array location (7) if the variable
is an array, the routine or common block in which the
variable is located (8), the internal representation
(9), and the decimal representation (10) of the value
contained in the location of the variable are reported.
The operand registers (11) contain the values which
were currently being used in the computation. The

decimal equivalents (12) of the contents of the operand registers are also supplied. The increment registers (13) usually contain counters such as the indices of DO loops. Their contents are useful in debugging simulations only very rarely. In (8) common blocks are represented by names enclosed in slashes. The blank common block whose entries are denoted by / / refers to the location of storage of reserved-system variables and state variables in addition to variables declared in storage-allocation statements. For the example simulation we find that Q contains an indefinite value. Referring back to the listing in Fig. 2.3-1, we find that Q was never assigned a value. This was the cause of the error. In subroutine PVAL we had attempted to add the value of Q to the quantity TIME/R, but Q had not been given a value. A special note of caution is in order. The failure to determine the cause of an error through the use of the information contained in the exchange jump package is usually caused by trying to digest too much information. Many times much of the information is not relevant to the discovery of the error. In determining the cause of the error in the above example, we had combined the information that the error was detected in subroutine PVAL *along* with the information that Q was indefinite. We were thus not mislead by the fact that

FLOW is indefinite. FLOW is indefinite because it is
not assigned a value until the call to PVAL is completed.

The third part of a debug report contains a listing
of the values of the variables when the error was detected.
Most of the information is self-explanatory. If a vari-
able is an array which contains successive equal values,
a repetition factor is used to conserve space (14). In
this example the first nine state variables are
indefinite. State variables X(10), X(11), and X(12)
contain respectively the values 100.0, indefinite and
0. The remaining 987 state variables are indefinite.
This is all satisfactory since X(10) and X(12) are the
only state variables used in this simulation and are
defined (i.e., given legal values). The variable XMFL
in the routine XFLOWS has a special meaning. If an
error is detected in routine XFLOWS, the value of XMFL
+ 1 points to the flow which was being computed at the
time of the error. If XMFL equals zero, the first flow
was being computed. If XMFL equals 10, the 11th flow
defined in the simulation was being computed. Do not
forget to count all flows in iteratively defined flows.

The following sample simulations are shown to
illustrate various types of errors and the procedure for
determining the cause of the errors using the debug
report as a guide. Each table following a listing and
a debug report contains the relevant information used

in deducing the cause of the error. Practice is required
in recognizing the relevant pieces of information which,
when combined, produces an explanation of the error. Many
times a single piece of information is misleading unless
it is interpreted along with other pieces of information.

Example 2.3-2.   Illustration of an uninitialized state variable (see Table
2.3-4).

```
DEBUG.
REAL. N
(2-3). FLOW=N*COS(TIME*6.28/50.)  (4)
7 8 9      end-of-record separator
TSTRT=0. $ TEND=100. $ DT=1. $ X(2)=100. $ N=10. $  (5)
PLOT. (X(2).X(3))
```

ARITHMETIC MODE ERROR   DIAGNOSTIC DUMP   07/23/73   10.51.57.

TYPE OF ERROR:
    ERROR MODE = 4   (1)

        FLOATING POINT ARITHMETIC UNIT RECIEVED AN INDEFINITE OPERAND

        OCCURING (APPROXIMATELY) AT ADDRESS 010667B WHICH IS LOCATION 000021B IN ROUTINE XCSIM (2)

        NON-STANDARD FLOATING POINT ARITHMETIC - TABLES OF NON-STANDARD RESULTS BY DIVISION

                    DIVIDE (A/B)   WHERE

                          B   +0   =   0000 X...X B
                    +N   -N   +0   -0   -0   =   7777 X...X B
                                                    +INF =   3777 X...X B   (+ INFINITY)
              +N   --   --   +INF -INF   -INF =   4000 X...X B   (- INFINITY)
          A   -N   --   --   -INF +INF   +IND =   1777 X...X B   (+ INDEFINITE)
              +0   0   0   +IND +IND   -IND =   6000 X...X B   (- INDEFINITE)
              -0   0   0   +IND +IND   N   =   ANY WORD EXCEPT +INF, -INF, +IND, -IND, +0, OR -0

    EXCHANGE JUMP PACKAGE:

    ADDRESS   REFERENCED   TYPE   INDEX   LOCAL   CONTAINED
    REGISTERS   CONTENTS   VARIABLE NAME **ARRAY (DEC) ADDRESS   IN   VALUE (OCTAL)   DECODED VALUE

        A0   052000R   OUT OF RANGE
        A1   014726B   DT   REAL   000004B   / /   1720 4000 0000 0000 0000B   1.000000000000
        A2   010705R   000037H   XCSIM   0000 0000 0077 7770 0000B   1073709056
        A3   010711B   000043B   XCSIM   5170 0313 4321 1444 6000B   -1.3490593968519+132
        A4   010706B   00004BR   XCSIM   0000 0000 0000 0007 7777B   32767
        A5   014734B   X   (3)   *REAL   3   000010R   / /   1777 0000 0000 0000 0013B   + INDEFINITE
        A6   014733R   X   *REAL   2   000010B   / /   1726 5500 0000 0000 0000B   90.00000000000
        A7   006020B   000000B   XXFL2WS 1723 4777 7777 7777 7777B   10.000000000000

    OPERAND   INCREMENT
    REGISTERS   CONTENTS   DECODED VALUE   REGISTERS   CONTENTS

        X0   1723 4777 7777 7777 7777R   10.000000000000   B0   0000008
        X1   1720 4000 0000 0000 0000R   1.000000000000   B1   0060208
        X2   0000 0000 0077 7770 0000R   1073709056   B2   0060178
        X3   0000 0000 0000 0000 0002R   2   B3   0107128
        X4   1726 5500 0000 0000 0000R   90.00000000000   B4   7777728
        X5   1777 0000 C000 0000 0013R   + INDEFINITE   B5   0000018
        X6   1777 0000 0000 0000 0000R   + INDEFINITE   B6   7777738
        X7   0000 0000 0000 0000 0003R   3   B7   0060178

```
VARIABLE DUMP  -  XFLOWS

   VARIABLE      TYPE                  LOCAL
     NAME      **APRAY   LOCATION    ADDRESS    REPLATED         VALUE (OCTAL)              DECODED VALUE

     XMFL      INTEGER   0076068    0000178                0040 0000 0090 0000 00018                          I
     FLOW      REAL      0076078    000020R                1723 4777 7777 7777 77778       10.000000000096


VARIABLE DUMP  -  / /

     XADRS     *REAL     0147228    0000008                5663 0663 3606 5301 34548       -1.9006556451889E+37
     TIME      REAL      0147238    000001R                0000 0000 0000 0000 00008       0.
     TSTRT     REAL      0147248    0000028                0000 0000 0000 0000 0000H       0.
     TEND      REAL      0147258    0000038                1776 6200 0000 0000 00008       100.0000000000
     DT        REAL      0147268    0000048                1720 4000 0000 0000 00008       1.000000000000
     DTPR      REAL      0147278    0000058                1777 0000 0000 0000 00008       + INDEFINITE
     DTPL      REAL      0147308    0000068                1720 4000 0000 0000 00008       1.000000000000
     DTFL      REAL      0147318    0000078                1777 0000 0000 0000 00008       + INDEFINITE
     X         *REAL     0147328    0000108                1777 0000 0000 0000 00008       + INDEFINITE
                                                           1726 5500 0000 0000 00008       90.00000000000
                                               997 _       1777 0000 0000 0000 00008       + INDEFINITE
     N         REAL      0167018    0017578                1723 5000 0000 0000 00008       10.00000000000
     XEVSTR    *REAL     0167028    0017608                0070 0000 0010 0001 04108       1.0768146479358-283
```

Table 2.3-4.   Information used in determining the error in example 2.3-2. Item numbers refer to the circled numbers in example 2.3-2.

Item No.	Information
(1) and (2)	An indefinite quantity was used in XCSIM. Therefore either a flow or state variable was indefinite.
(3)	X(3) was indefinite.
(4) and (5)	X(3) is the source compartment for the flow, but was not initialized in the data section.

Example 2.3-3.   Illustration of the failure to assign a value to FLOW.

(see Table 2.3-5).

```
DEBUG.
(1-2). V=X(1)*0.01
       W=X(2)*COS(TIME*6.28/50.)     ⑥
       FLW=W+V
?8ç      end-of-record separator

TSTRT=0. $ TEND=100. $ DT=1. $ X=100..1. $
PLOT. (X(1)),(X(2))
```

ARITHMETIC MODE ERROR   DIAGNOSTIC DUMP   07/23/73   20.35.34.

TYPE OF ERROR:
     ERROR MODE = 4   ①

          FLOATING POINT ARITHMETIC UNIT RECIEVED AN INDEFINITE OPERAND

     OCCURING (APPROXIMATELY) AT ADDRESS 0106710 WHICH IS LOCATION 4000160 IN ROUTINE XCSIM ②

     NON-STANDARD FLOATING POINT ARITHMETIC - TABLES OF NON-STANDARD RESULTS BY DIVISION

                         DIVIDE (A/B)   WHERE

                              B   +0   =   0000 X...X B
                        +N   -N   +0   -0   -0   =   7777 X...X B
                                                        +INF =   3777 X...X B (+ INFINITY)
                  +N   --   --   +INF -INF   -INF =   4000 X...X B (- INFINITY)
              A   -N   --   --   -INF +INF   +IND =   1777 X...X B (+ INDEFINITE)
                  +0   0   0   +IND +IND   -IND =   0000 X...X B (- INDEFINITE)
                  -0   0   0   +IND +IND   N   =   ANY WORD EXCEPT +INF, -INF, +IND, -IND, +0, OR -0

EXCHANGE JUMP PACKAGE:

ADDRESS REGISTERS	CONTENTS	REFERENCED VARIABLE NAME	TYPE **=ARRAY	INDEX (DEC)	LOCAL ADDRESS	CONT/INED IN	VALUE (OCTAL)	DECODED VALUE	
A0	0520008						OUT OF RANGE		
A1	0147330	DT	REAL	0000048	/ /	1720 4000 0000 0000 00008	1.000000000000	
A2	0107128			000037R	XCSIM	0000 0000 0077 7770 00008	1073709056	
A3	0107160			0000438	XCSIM	5957 3342 5555 5733 4055B	-3.575749942679TE+57	
A4	0060100			0000000	XXFL2WS	1777 0000 0000 0000 00018	+ INDEFINITE	
A5	0060150			0000010	XXFL1WS	0000 0000 0000 0010 00028	32770	
A6	0060160			000000R	XXFL2WS	1777 0000 0000 0000 00018	+ INDEFINITE	
A7	0076100	XMFL	INTEGER	0000230	XFLO>S	0000 0000 0000 0000 00018	1	1

OPERAND REGISTERS	CONTENTS	DECODED VALUE		INCREMENT REGISTERS	CONTENTS
X0	1777 0000 0000 0000 0000R	+ INDEFINITE		B0	0000008
X1	1720 4000 0000 0000 0000R	1.000000000000		B1	0060168
X2	0000 0000 0077 7770 0000R	1073709056		B2	0060158
X3	0000 0000 0000 0000 0001R	1		B3	0107178
X4	1777 0000 0000 0000 00018	+ INDEFINITE		B4	7777728
X5	0000 0000 0000 0010 0002R	32770		B5	0000010
X6	1777 0000 0000 0000 0001R	+ INDEFINITE		B6	7777738
X7	0000 0000 0000 0000 0001R	1		B7	0060158

VARIABLE DUMP - AFLOWS

VARIABLE NAME	TYPE **=ARRAY	LOCATION	LOCAL ADDRESS	REPEATED	VALUE (OCTAL)	DECODED VALUE	
XMFL	INTEGER	0076105	000023R		0000 0000 0000 0000 00018	+ INDEFINITE	1 ③
FLOW	REAL	0076118	0000248		1777 0000 0000 0000 00008	1.000000000000	
V	REAL	0076128	0R0025R		1720 4000 0000 0000 00008	1.0000000000000	
W	REAL	0076138	0000268		1717 7777 7777 7777 77778	1.0000000000000	
FLW	REAL	0076148	0000278		1721 4000 0000 0000 00008	2.000000000000	⑤

VARIABLE DUMP - / /

XADRS	*REAL	0147277	0000008		5663 0663 3606 5301 34618	-1.9006556451889E+37	
TIME	REAL	0147308	000001R		0000 0000 0000 0000 00008	0.	
TSTRT	REAL	0147318	000002R		0000 0000 0000 0000 00008	0.	
TEND	REAL	0147328	000003R		1726 6200 0000 0000 00008	100.0000000000	
DT	REAL	0147338	0000048		1720 4000 0000 0000 00008	1.000000000000	
DTPR	REAL	0147348	0000058		1777 0000 0000 0000 00008	+ INDEFINITE	
DTPL	REAL	0147358	0000060		1720 4000 0000 0000 00008	1.000000000000	
DTFL	REAL	0147368	0000078		1777 0000 0000 0000 00008	+ INDEFINITE	
X	*REAL	0147378	000010R		1726 6200 0000 0000 00008	100.0000000000	④
						1720 4000 0000 0000 00008	1.000000000000	
					997	1777 0000 0000 0000 00008	+ INDEFINITE	
XEVSTK	*REAL	0167008	0017578		0100 0000 0010 0001 04158	2.7565976013290-281	

Table 2.3-5.   Information used in determining the error in example 2.3-3.

Item No.	Information
(1) and (2)	An indefinite quantity was used in XCSIM. Therefore either a flow or a state variable was indefinite.
(3)	The value of FLOW is indefinite.
(4)	The values of the state variables X(1) and X(2) are legal values.
(5) and (6)	A variable FLW contains a legal value but is obviously a mispunch for the variable FLOW.

Example 2.3-4.   Illustration of the generation of an infinite operand (see Table 2.3-6).

```
DEBUG.
STORAGE. BIRT,POP,DENS,AREA,ARMIN
        EVENT POPL
        CALL EVENT(4HPOPL,TIME+1.,1)
        BIRT=3.*SIN(TIME*6.28/50.)
        POP=POP+BIRT/DENS
        DENS=POP/AREA          ⑤
        RETURN
        END
        EVENT CRWD
        CALL EVENT(4HCRWD,TIME+1.,1)
        AREA=EXP(-TIME/10.)
        IF(AREA.LT.0.5) AREA=ARMIN   ⑧
        RETURN
        END
   7⁸₉    end-of-record separator

EVENT. POPL,0.
EVENT. CRWD,0.     ④   ⑥
EVENT. HALT,100.
PCP=3. $ DENS=3. $ AREA=1. $ ARMIN=0. $
PLOT. (BIRT),(POP),(DENS)
PLOT. (POP)/DENS              ⑩
```

ARITHMETIC MODE ERROR   DIAGNOSTIC DUMP   07/23/73   18.57.46.

TYPE OF ERROR:
    ERROR MODE = 2   ①

        FLOATING POINT ARITHMETIC UNIT RECIEVED AN INFINITE OPERAND

    OCCURING (APPROXIMATELY) AT ADDRESS 0076368 WHICH IS LOCATION 0000128 IN ROUTINE POPL ②

    NON-STANDARD FLOATING POINT ARITHMETIC - TABLES OF NON-STANDARD RESULTS BY DIVISION

                    DIVIDE (A/B)   WHERE

                            B   +0   =   0000 X...X B
                                                     -0   =   7777 X...X B
                    +N   -N   +0   -0   +INF =   3777 X...X B (+ INFINITY)
                                                     -INF =   4000 X...X B (- INFINITY)
            +N   --   --   +INF -INF   +IND =   1777 X...X B (+ INDEFINITE)
        A   -N   --   --   -INF +INF   -IND =   6000 X...X B (- INDEFINITE)
            +0   0   0   +IND +IND   N   =   ANY WORD EXCEPT +INF, -INF, +IND, -IND, +0, OR -0
            -0   0   0   +IND +IND

EXCHANGE JUMP PACKAGE:

ADDRESS REGISTERS	CONTENTS	REFERENCED VARIABLE NAME	TYPE *=ARRAY	INDEX (DEC)	LOCAL ADDRESS	CONTAINED IN	VALUE (OCTAL)	DECODED VALUE
							OUT OF RANGE	
A0	0520008							
A1	0134358				0000408	SINCOSE	6185 3301 0145 2401 66178	-2.7555218727710E-07
A2	0134368				0000418	SINCOSE	1663 4334 1433 4416 36078	2.0629106347665E-09
A3	0134418				0000448	SINCOSE	1713 5252 5252 5252 34678	4.1666666666670E-02
A4	0134428				0000458	SINCOSE	6061 0000 0000 0000 00048	-.5000000000000
A5	0170468	DENS	REAL		0017618	//	3777 0000 0000 0000 00008	+ INFINITE ③
A6	0056718				0000028	XXEVENT	0000 0000 0000 0000 00108	8
A7	0076468				0000228	POPL	1723 5000 0000 0000 00008	10.00000000000

OPERAND REGISTERS	CONTENTS	DECODED VALUE		INCREMENT REGISTERS	CONTENTS
X0	0000 0000 0000 0000 00008	0		B0	0000008
X1	6073 3454 2514 6352 6171R	-2.69373506094458E-04		B1	0000018
X2	1673 7636 2501 0067 0417R	9.30933368490074E-07		B2	0000018
X3	0000 0000 0000 0000 00008	0		B3	0000008
X4	6061 0000 0000 0000 0004R	-.5000000000000		B4	0000578
X5	3777 0000 0000 0000 0000R	+ INFINITE		B5	0000038
X6	1717 7171 1274 5512 71448	.9845827809445		B6	0170578
X7	1721 5332 7015 4170 1313R	2.7137483+2833		B7	0000068

VARIABLE DUMP - XFLOVS

VARIABLE NAME	TYPE *=ARRAY	LOCATION	LOCAL ADDRESS	REPEATED	VALUE (OCTAL)	DECODED VALUE
XMFL	INTEGER	007623R	0000048		0003 5003 1700 0000 00008	X

VARIABLE DUMP - //

XADRS	*REAL	0150658	000000R		5663 0663 3606 5301 36178	-1.9806556451881E+37
TIME	REAL	0150668	000001R		1723 4400 0000 0000 00008	9.000000000000
TSTRT	REAL	0150678	0000028		0000 0000 0000 0000 00008	0.
TEND	REAL	0150708	0000038		1777 0000 0000 0000 00008	+ INDEFINITE
DT	REAL	0150718	0000048		1777 0000 0000 0000 00008	+ INDEFINITE
DTPR	REAL	0150728	0000058		1777 0000 0000 0000 00008	+ INDEFINITE
DTPL	REAL	0150738	0000068		1720 4000 0000 0000 00008	1.000000000000
DTFL	REAL	0150748	0000078		1777 0000 0000 0000 00008	+ INDEFINITE
X	*REAL	0150758	0000108	999	1777 0000 0000 0000 00008	+ INDEFINITE
RIRT	RFAL	0170448	0017578		1721 5040 7372 3261 37048	2.532104196570
POP	REAL	0170458	0017608		1722 5156 1066 0240 13508	5.215113658479
DENS	REAL	0170468	0017618		3777 0000 0000 0000 00008	+ INFINITE ⑦
AREA	REAL	0170478	0017628		0000 0000 0000 0000 00008	0.
AMMIN	REAL	0170508	0017638		0000 0000 0000 0000 00008	0.
REVSTK	*REAL	0170518	0017648		0104 0000 0000 0101 05078	0.7530724920936-283 ⑧

Table 2.3-6.   Information used in determining the error in example 2.3-4.

Item No.	Information
(1) and (2)	An infinite operand was used in event POPL.
(3)	The variable DENS was + infinite.
(4)	DENS was initialized to a legal value.
(5)	DENS is computed as POP/AREA.
(6) and (7)	AREA was initialized to 1., but now contains the value 0.
(8)	AREA is recomputed in event CRWD and can assume the value of ARMIN.
(9) and (10)	ARMIN currently has the value zero and was mistakenly initialized to zero.

Example 2.3-5.   Illustration of an out-of-range subscript (see Table 2.3-7).

```
DEBUG.
STORAGE.  P(5).N          (5)
(1-I=11,15).  J=I+N
      FLOW=P(J)*X(I)

78g      end-of-record separator                          (8)

TSTRT=0.  $ TEND=10.  $ DT=1.  $ DTPR=1.  $
P=3*0.1,0.2,0.5 $ X(1)=1000.  $ X(11)=30.,20.,25.,2*10.  $ N=9000 $
PRINT.
```

ARITHMETIC MODE ERROR   DIAGNOSTIC DUMP   07/23/73 20.06.33.

TYPE OF ERROR:
 ERROR MODE = 1   (1)

  ATTEMPTED TO REFERENCE CENTRAL MEMORY OUTSIDE ESTABLISHED LIMITS

  OCCURING (APPROXIMATELY) AT ADDRESS 007614B WHICH IS LOCATION 0000138 IN ROUTINE XFLOWS (2)

EXCHANGE JUMP PACKAGE:

ADDRESS REGISTERS	CONTENTS	REFERENCED VARIABLE NAME	TYPE *=ARRAY	INDEX (DEC)	LOCAL ADDRESS	CONTAINED IN	VALUE (OCTAL)	DECODED VALUE
A0	052000B						OUT OF RANGE	
A1	010724B				000041B	XCSIM	0000 0000 0000 0000 0000B	0
A2	040323B						OUT OF RANGE	
A3	014704B	X	*REAL	11	000010B	/ /	1724 7400 0000 0000 0000B	30.00000000000
A4	007624B	J	INTEGER		0000230	XFLOWS	0000 0000 0000 0002 1463B	9011
A5	007622B	I	INTEGER		000021B	XFLOWS	0000 0000 0000 0000 0013B	11
A6	007622B	I	INTEGER		000021B	XFLOWS	0000 0000 0000 0000 0013B	11
A7	007624B	J	INTEGER		000023B	XFLOWS	0000 0000 0000 0002 1463B	9011

                                 (4)

OPERAND REGISTERS	CONTENTS	DECODED VALUE	INCREMENT REGISTERS	CONTENTS
X0	0000 0000 0000 0000 0002R	2	B0	000000B
X1	0000 0000 0000 0000 0000R	0	B1	014704B
X2	0000 0000 0000 0000 0000R	0	B2	000001B
X3	1724 7400 0000 0000 0000R	30.00000000000	B3	007624B
X4	0000 0000 0000 0002 1463R	9011	B4	000004B
X5	0000 0000 0000 0000 0013R	11	B5	000063B
X6	0000 0000 0000 0000 0013R	11	B6	000001B
X7	0000 0000 0000 0002 1463R	9011	B7	014710B

VARIABLE DUMP - XFLOWS

VARIABLE NAME	TYPE *=ARRAY	LOCATION	LOCAL ADDRESS	REPEATED	VALUE (OCTAL)	DECODED VALUE
XMFL	INTEGER	007621B	000020R		0000 0000 0000 0000 0000B	0
I	INTEGER	007622B	000021B		0000 0000 0000 0000 0013B	11
FLOW	REAL	007623B	000022R		0003 5003 1700 0000 0000B	1.567734751+039-293
J	INTEGER	007624R	000023B		0000 0000 0000 0002 1463B	9011

VARIABLE DUMP - / /

XAORS	*REAL	014662B	000000R		5663 0663 3606 5301 3414B	-1.900655645189IE+37
TIME	REAL	014663B	000001R		0000 0000 0000 0000 0000B	0.
TSTRT	REAL	014664B	0000028		0000 0000 0000 0000 0000B	0.
TEND	REAL	014665R	0000038		1723 5000 0000 0000 0000B	10.00000000000
D:	REAL	014666B	0000048		1720 4000 0000 0000 0000B	1.000000000000
DTPR	REAL	014667B	0000058		1720 4000 0000 0000 0000B	1.000000000000
DTPL	REAL	014670R	0000068		1777 0000 0000 0000 0000B	+ INDEFINITE
DTFL	REAL	014671B	000007R		1777 0000 0000 0000 0000B	+ INDEFINITE
X	*REAL	014672B	000010B		1731 7640 0000 0000 0000B	1800.000000000
				9 _	1777 0000 0000 0000 0000R	+ INDEFINITE
					1724 7400 0000 0000 0000B	30.00000000000
					1724 5000 0000 0000 0000B	20.00000000000
					1724 6200 0000 0000 0000B	25.00000000000
				2 _	1723 5000 0000 0000 0000B	10.00000000000
				984 _	1777 0000 0000 0000 0000B	+ INDEFINITE
P	*REAL	016641B	0017578	3 _	17.4 6314 6314 6314 6315B	.1000000000000
					1715 6314 6314 6314 6315B	.2000000000000
					1717 4000 0000 0000 0000B	.5000000000000
N	INTEGER	016646B	0017648		0000 0000 0000 0002 1450B	9000
XEVSTK	*REAL	016647B	0017658		0074 0000 0010 0001 0425B	1.721623603444D-282

Table 2.3-7. Information used in determining the error in example 2.3-5.

Item No.	Information
(1) and (2)	An attempt was made to reference central memory outside established limits in routine XFLOWS. Therefore the error occurred while evaluating a flow.
(3)	XMFL equals zero; therefore the first flow was being computed.
(4)	The value of J is 9011; much too large.
(5)	J is defined as the sum of I and N.
(6)	N was initialized to 9000, an error. N should have been initialized to -10.

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APPENDIX A

DIAGNOSTICS

Compilation Diagnostics

The SIMCOMP compiler will list error messages immediately following

the source card containing the rule infraction. Errors are either fatal,

which are prefaced by *****FE, or nonfatal, which are prefaced by *****NF.

Fatal errors cause abnormal termination, and execution of the simulation

will not occur. Nonfatal errors only result in a diagnostic being issued,

but should be corrected in a subsequent run. Unpredictable results in the

execution are possible if nonfatal errors exist in the source section.

```
*****FE    A FIELD IN WHICH A CONSTANT SHOULD APPEAR IS MISSING OR IS NEGATIVE
*****FE    A VARIABLE DECLARATION IS INCOMPLETE AT CARD END
*****FE    ABOVE CARD ILLEGAL AT THIS POINT
*****FE    ARITHMETIC PHRASE MUST BE USED IN CONJUNCTION WITH A DO... PHRASE
*****FE    CHARACTER "_" IS ILLEGAL
*****FE    CHARACTER "_" IS ILLEGAL IN COLUMN __
*****FE    EXPECTED SUBSCRIPT MISSING
*****FE    EXPECTED VARIABLE NAME MISSING
*****FE    FLOW DIRECTIVE UNTERMINATED AT CARD END
*****FE    FLOW EXPRESSION SUB-FIELD "_____..." CONTAINS MORE THAN 10 NON-BLANK CHARS
*****FE    FLOW INDICES ( ___-___ ) PRODUCED BY THE ABOVE LABEL ARE OUTSIDE THE RANGE 1 - 999
*****FE    FLOW ITERATION PHRASE CONTROL VARIABLE "_____" MUST BE A 5 CHAR OR LESS INTEGER VARIABLE
*****FE    FLOW PHRASE "_____..." CONTAINS MORE THAN 40 NON-BLANK CHARS
*****FE    INSUFFICIENT FIELD LENGTH, INCREASE BY (NO. OF FLOWS - ___)
*****FE    NUMBER OF DECLARED VARIABLES HAS EXCEEDED _____
*****FE    NUMBER OF FLOWS EXCEEDS 9999
*****FE    NUMBER OF USER-DEFINED EVENTS EXCEEDS 100
*****FE    ROUTINE NAME LONGER THAN 5 CHARS OR MISSING
*****FE    ROUTINE NAME MISSING
*****FE    ROUTINE NAME "_____..." LONGER THAN 7 CHARS
*****FE    ROUTINE NAME "_____' STARTS WITH AN ILLEGAL CHAR
*****FE    SUBSCRIPT "_____..." IS LONGER THAN 4 CHARS
*****FE    SUBSCRIPT "_____" GREATER THAN 1023
*****FE    SUBSCRIPT "_____" NOT DECODABLE
*****FE    THE DO... PHRASE CONTROL VARIABLE MUST BE THE OPERAND IN THE ARITHMETIC PHRASE
*****FE    VARIABLE "_____..." IS LONGER THAN 5 CHARS
*****FE    VARIABLE "_____" BEGINS WITH A NON-ALPHABETICAL CHAR
*****NF    VARIABLE "_____" BEGINS WITH CHAR "X"
*****NF    VARIABLE "_____" HAS BEEN PREVIOUSLY DECLARED, LAST DECLARATION IS ASSUMED CORRECT
*****FE    VARIABLE "_____" IS A RESERVED SYSTEM VARIABLE
```

Data Section Diagnostics

Errors encountered while processing the data section are reported in the output by a general message of the following form:

```
***** ERROR IN PRINT REQUEST
***** ERROR IN FLOW PRINT REQUEST
***** ERROR IN EXOGENOUS EVENT REQUEST
***** ERROR IN DATA ASSIGNMENT
***** ERROR IN PLOT REQUEST
```

One of these messages is followed by the card containing the infraction and one of the following diagnostics. All errors reported in the data section are nonfatal. The system will attempt to execute the simulation regardless of errors in the data section. Execution errors can occur because of errors in data assignment and exogenous-event requests.

```
CHARACTER "_" IS ILLEGAL IN COLUMN __
DATA ITEM "_____..." LONGER THAN 20 CHARS
DATA REPETITION FACTOR "_____..." LONGER THAN 10 CHARS
DATA REPETITION FACTOR "_____" LESS THAN OR EQUAL TO ZERO
DATA REPETITION FACTOR "_____" NOT DECODABLE
EVENT NAME "_____..." LONGER THAN 5 CHARS
EVENT "_____" IS NON-EXISTANT
EVENT "_____" SCHEDULED AT TIME _____ AT PRIORITY OF __
EXPECTED FLOW INDEX MISSING IN OR BEFORE COLUMN __
EXPECTED VARIABLE NAME MISSING IN OR BEFORE COLUMN __
FLOW INDEX "___.·.." LONGER THAN 3 CHARS
FLOW INDEX "___" NOT DECODABLE OR OUT OF RANGE
FLOW INDICES UNTERMINATED AT CARD END
FLOW PRINTING REQUESTED - NO FLOWS DEFINED
FLOW (__,__) DOES NOT EXIST
ILLEGAL CHARACTER DETECTED "_"
ILLEGAL CHARACTER IN RANGE DECLARATION
IMPROPERLY FORMATTED LOG REQUEST
INTEGER VARIABLE _____ WAS ASSIGNED A REAL VALUE IN THE DATA SECTION
MISSING EXPECTED VARIABLE NAME OR DATA ITEM IN OR BEFORE COLUMN
MORE THAN 100 VARIABLES NAMED IN PLOT REQUESTS, THIS AND SUBSEQUENT PLOT REQUESTS IGNORED
MORE THAN 200 VARIABLES REQUESTED FOR PRINT
NO. OF EXOGENOUSLY SCHEDULED EVENTS EXCEEDS 20, ABOVE REQUEST IGNORED
NO. OF GROUPS PER PLOT IS .GT. 5
NO. OF VARIABLES PER PLOT IS .GT. 5
RANGE DECLARATION .GT. 10 CHARACTERS--THE UPPER LIMIT
REAL VARIABLE _____ WAS ASSIGNED AN INTEGER VALUE IN THE DATA SECTION .
SUBSCRIPT "___..." LONGER THAN 4 CHARS IN COLUMN __
SUBSCRIPT "___" NOT DECODABLE
TIME OR PRIORITY LONGER THAN 20 CHARS AT COLUMN __
VARIABLE HAS .GT. 3 SUBSCRIPTS
VARIABLE NAME IS .GT. 5 CHARACTERS
VARIABLE SUBSCRIPT .GT. 999--THE UPPER LIMIT
VARIABLE "_____..." LONGER THAN 5 CHARS
VARIABLE "_____" WAS NOT COMPLETELY DECLARED BY CARD END
VARIABLE "_____" WAS NOT DECLARED IN A <STORAGE.> STATEMENT
```

APPENDIX B

DECK ORGANIZATION AND CONTROL CARDS

A typical SIMCOMP job is executed by means of the following control
cards.

```
TAxxx, Annnnnnn.  (job card)
ATTACH,SIMCOM,SIMCOM3,CY=1,MR=1,ID=NREL.
SIMCOM.
```

$^7 8_9$ (end of record)

       .
       .
       .

    source section

       .
       .

$^7 8_9$

       .
       .

    data section

       .
       .

$^6 7 8_9$ (end of file)

In actuality more control cards than those shown are utilized in executing the simulation. The SIMCOMP compiler generates a series of control cards which are used subsequent to the loading and execution of the compiler. As described in section 2.1.7, a SIMCOMP simulation can be executed in three different modes. If a NOGO. execution directive is included in the source section, the generation of these control cards by the compiler is inhibited. If the default mode or the DEBUG. mode of execution is selected, standard sets of control cards are generated automatically which are used after the SIMCOM. control card is executed. If a NOGO. directive is included in the source section, the following control cards and deck structure is equivalent to the deck in the above example with the default execution mode selected,

i.e., by the absence of any execution directives. In the following case the

user is supplying the control cards rather than having the compiler generate

them automatically.

```
TAxxx, Annnnnnn.  (job card)
ATTACH,SIMCOM,SIMCOM3,CY=1,MR=1,ID=NREL.
SIMCOM.
FTN,I=SIMPRG,ROUND=T-*/,S=0,LRN=0.
ATTACH,B,SIMCOM3,CY=2,MR=1,ID=NREL.
ATTACH,LIB,SIMCOM3,CY=3,MR=1,ID=NREL.
SELECT.
COPYBF,B,LGO.
LOAD,LGO.
NOGO.1/
REWIND,NEWT1.
SELECT,P=PRELOAD,I=PRELOAD.
PRELOAD,NEWT1,MAIN.
MAIN.
```

$^7 8_9$

NOGO.

⋮

source section

⋮

$^7 8_9$

⋮

data section

⋮

$^6 {}^7 8_9$

Similarly the following example is equivalent to the first example if

a DEBUG. directive had been included in the source section. Here again the

user is supplying the required additional control cards since the automatic

_____

[1]/ Not to be confused with the special execution directive NOGO. which
    is included in the source section.

generation of the control cards is suppressed by the NOGO. directive in the source section.

```
TAxxx, Annnnnnn.   (job card)
ATTACH,SIMCOM,SIMCOM3,CY=1,MR=1,ID=NREL.
SIMCOM.
FTN,I=SIMPRG,LN=DEBUG,R=1,S=0,ROUND=T-*/.
ATTACH,B,SIMCOM3,CY=2,MR=1,ID=NREL.
ATTACH,LIB,SIMCOM3,CY=3,MR=1,ID=NREL.
SELECT.
COPYBF,B,LGO.
MAP,PART.
LOAD,LGO.
NOGO.
REWIND,NEWT1.
SELECT,P=PRELOAD,I=PRELOAD.
PRELOAD,NEWT1,MAIN.
MAIN.
```

$^7 8_9$

NOGO.

$\vdots$

source section

$\vdots$

$^7 8_9$

$\vdots$

data section

$\vdots$

$^6 {}^7 8_9$

Job Limits. The job card illustrated in the above examples implicitly requests the minimum amount of time, pages printed, cards punched, and core required for a simple SIMCOMP job. These limits are:

Limit	Mnemonic on Job card	Meaning
Time	T16	16 seconds CPU time
Core	CM43000	43000 octal words of central memory
Printed pages	PR10	10 printed pages
Punched cards	PU10	10 punched cards

These limits are usually adequate only for the smaller SIMCOMP jobs. The limits specified on the job card should reflect the physical size of the simulation and the number of time steps (i.e., (TEND-TSTRT)/DT or the total number of events executed) used during the execution of the simulation. This is true for the time and core requirements. The number of printed pages is a function of the listing length and the number of output requests. A punch limit is required only if routine PUNCHD is called (refer to section 2.1.5). Only experience can be used to estimate these limits.

INDEX